### An Introduction to CUDA and GPU Programming

Scientific Computing and Data Analysis



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## GPU

- GPU is a "Graphics Processing Unit"
  - Developed for games
  - High computational capability



- CUDA is NVIDIAs C/C++ language extension for HPC computing on GPUs
  - But other options exist: OpenCL, OpenACC...



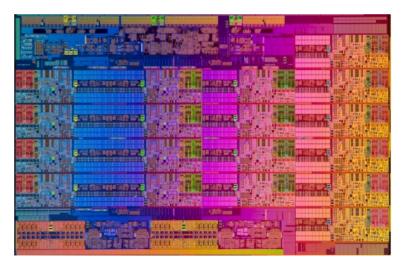
## Trend: multiple cores, parallel execution

As transistor count grows, CPUs have gained more cores and more features.

**A CPU** is a collection of 2-20 independent, fully **general cores** connected in a local network, with local and shared memory, and I/O to the outside world.

Xeon E5 2699v3

5.5B transistors18 cores45 MB L3 cache



## Meanwhile...

- PC graphics gained hardware support, 2D then 3D
- Early 2000s: GPUs with on-card programmable 3D transformation and lighting calculations ("shaders")

A shader is a *small* program that:

- Runs on many pixels or many vertices at once ("SIMD")
- Is compute bound only math with no I/O or conditionals
- Is stream oriented can process a stream of data points with little or no additional state



## Meanwhile...

- Researchers realized that you can run any function as a shader First only a curiosity, but it became obvious that GPUs were potentially powerful general code accelerators.
- 2007: NVIDIA releases CUDA, Apple releases OpenCL (later standardized under Chronos).

Both are direct means of harnessing the GPU compute power without hacks and workarounds.







#### CPU programs are **complex**:

• OS, simulators, games...

#### Cores are independent:

• Each core is a full computer

#### General computing tasks:

• Math, I/O, interactive applications.

#### GPU programs are **simple**:

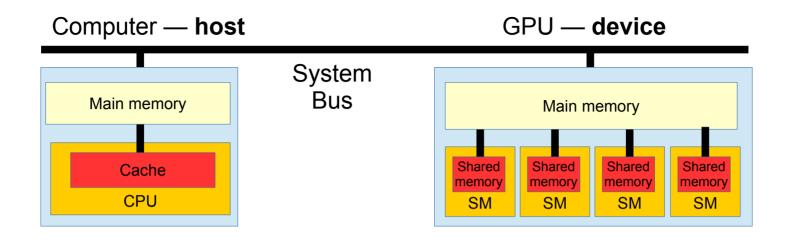
• Small size, simple control flow

#### Cores are very **parallel**:

 Many execution threads share one control unit

#### Very **specialized** for math:

• floating point function evaluators.



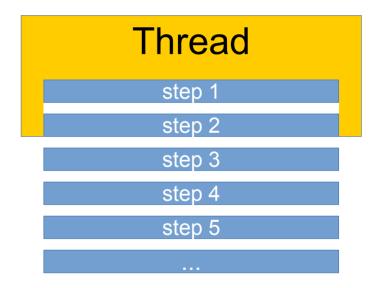
- The CPU ("host") and the GPU ("device") are separate
   → you need to copy data to and from the GPU
- The GPU has "streaming multiprocessors" = processors that each can run hundreds of threads.
  - Each SM has 8-64 CUDA cores; with more threads they take turns on the SM



## GPU concepts

A thread runs a single computation.

- Like CPU threads, they share memory and code with other threads.
- Much simpler, slower than CPU cores.
- Limited thread-local memory, registers.

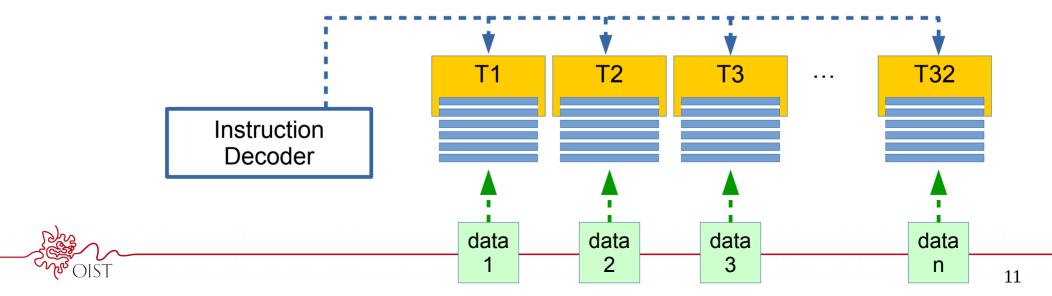




### A warp is a collection of 32 threads

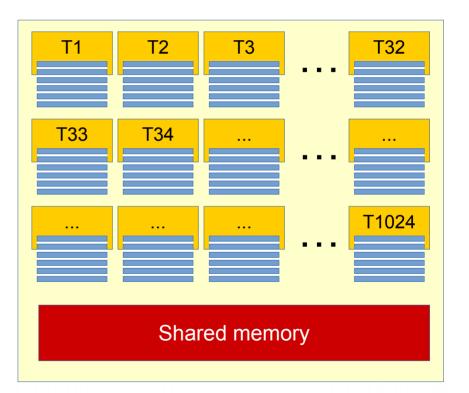
All threads in a warp run the same code at the same time
 → one thread takes less space and energy than a CPU core

- Threads have their own registers and variables.
- Same instructions, but act on different data ("SIMD")



## Blocks

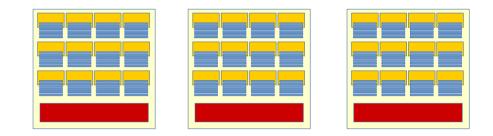
- A **block** is a collection of up to 1024 threads
  - All threads run the same code
  - Has shared fast memory (48KB)
- You can organize threads in 1-D,
   2-D or 3-D, but that's only
   programmer convenience



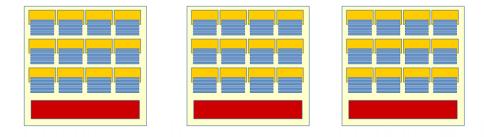
## Grid

#### A grid is a set of blocks

- Blocks are independent
  - Can not access data in shared memory of other blocks
- Blocks run in **any order.**
- Threads in all blocks have main memory in common
- You can organize blocks in different dimensions, the same as threads in a block.



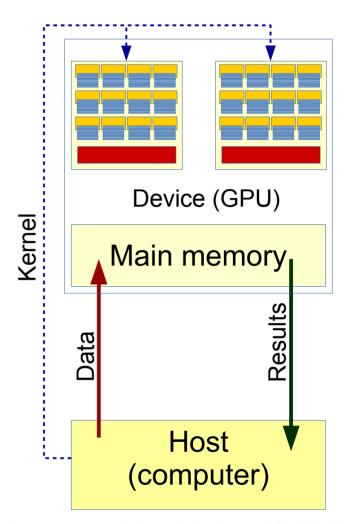
Main (device) memory





## Workflow

- 1) Write a C function (a "*kernel*") that will run on all **threads**
- 2) Copy your data from host to device
- Run the kernel on the device with your data, using some number of blocks and threads
- 4) Copy results from device back to host
- 5) Repeat from 3); or finish up
- 6) Done!





## Log In on Tombo!

• First, let's all log in on Tombo:

\$ ssh <your-ID>@tombo.oist.jp

• Copy the code from the common area:

\$ cp -r /work/training/GPU .

• Go to the new directory:

\$ cd GPU/code



### **GPU Resources at OIST**

#### Tombo:

"gpu"

- 1 node 2\* K40

#### Sango:

"gpu" - 2 node 4\* K80 "powernv" - 4 node 4\* P100

#### Saion:

"gpu"

- 9 node 4\* P100 - 8 node 4\* V100

"powernv" - 4 node 4\* V100

# \$ ssh <your-ID>@tombo.oist.jp \$ cp -r /work/share/training/GPU \$ ad CDU/aada

\$ cd GPU/code

### Training

image analysis pipelines general, deep learning

general, deep learning

general, deep learning

\$ ssh <your-ID>@tombo.oist.jp
\$ cp -r /work/share/training/GPU
\$ cd GPU/code

• You need to ask for the GPU partition, *and* reserve the GPU resource:

\$ srun --partition=gpu --gres=gpu <program>

- Use multiple cards with --gres=gpu:N
- For compilation, load the cuda module:

\$ module load cuda/8.0.27



### Program #1: GPU vector addition

We will adapt a simple program to run on the GPU:

```
void vec_add(float *a, float *b, float *c, int n) {
    int index;
    for (index=0; index<n; index++) {
        c[index] = a[index] + b[index];
    }
}</pre>
```

Vector addition is the "hello world" of parallel programming The source is GPU/code/vec\_add.c



The CPU-based version

- Allocate three arrays:
  - a, b inputs c output
- Initialize all elems
  - a,b = 5.0 c = 0.0
- Call vec\_add()

```
void vec_add(float *a, float *b, float *c, int n) {
    int index;
    for (index=0; index<n; index++) {</pre>
        c[index] = a[index] + b[index];
    }
}
int main(int argc, char **argv) {
    int i;
    float *c, *b, *a;
    c = (float *)malloc( size );
    b = (float *)malloc( size );
    a = (float *)malloc( size );
    for (i = 0; i<N; i++ ) {
        a[i] = b[i] = 5.0;
        c[i] = 0.0;
    }
    vec_add(a,b,c,N);
```

#### We have two extra things in our code. Timing:

```
double time_diff_nano(struct timespec *toc, struct timespec *tic) {
    return (1e9*(toc->tv_sec-tic->tv_sec)+
        (toc->tv_nsec-tic->tv_nsec));
}
...
clock_gettime(CLOCK_REALTIME, &tic);
...
clock_gettime(CLOCK_REALTIME, &toc);
time diff nano(&toc, &tic)/1000.0;
```

#### and sanity check:

```
float sum = 0.0;
for (index = 0; index<N; index++) {
    sum += c[index];
}
if (fabs((sum-CORRECT)/sum)>EPSILON) {
    printf(" correct sum: %.2f - fail\n", CORRECT);
}
```



### Run the serial code

• Compile our example program:

<pre>\$ gcc -03 vec_add.c -o addcpu</pre>	
Run it on Tombo:	On some systems you may need '-lrt' for the timer.

\$ srun -t 1:00 --mem=50m ./addcpu calculated sum: 10485760.00 Time: 1441.77 μs

• Time: ~1425-1550µs



### Allocate memory on the GPU

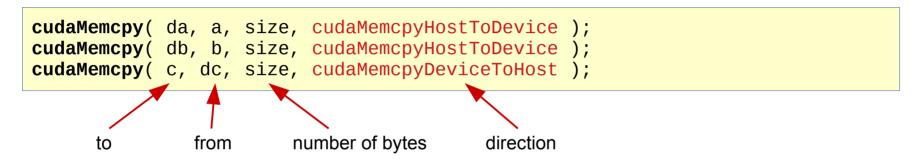
- Make a copy (or just use code/vec\_add.cu) and open it:
  \$ cp vec\_add.c vec\_add.cu
- Allocate memory on the GPU (and free it at the end):

```
float *dc, *db, *da;
cudaMalloc( &da, size );
cudaMalloc( &db, size );
cudaMalloc( &dc, size );
cudaFree( da );
cudaFree( db );
cudaFree( dc );
```

- All cuda functions start with "cuda"
- You usually don't need to add #include statements the cuda compiler adds then for you.

#### Move our data to the GPU

• Copy our data over:



cudaMemcpyHostToDevice = copy to GPU from computer cudaMemcpyDeviceToHost = copy to computer from GPU

also cudaMemcpyHostToHost, cudaMemcpyDeviceToDevice

Many function variations available: cudaMemcpyAsync(), cudaMemcpy2DTo[From]Array(), cudaMemcpyToSymbol() ...

#### Create the kernel

Our original vector addition function:

```
void vec_add(double *a, double *b, double *c, int n) {
    int index;
    for (index = 0; index<n; index++) {
        c[index] = a[index] + b[index];
    }
}</pre>
```

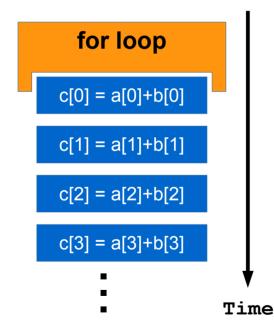
#### Add "\_\_global\_\_" specifier:

\_\_global\_\_ void vec\_add(double \*a, double \*b, double \*c, int n) {

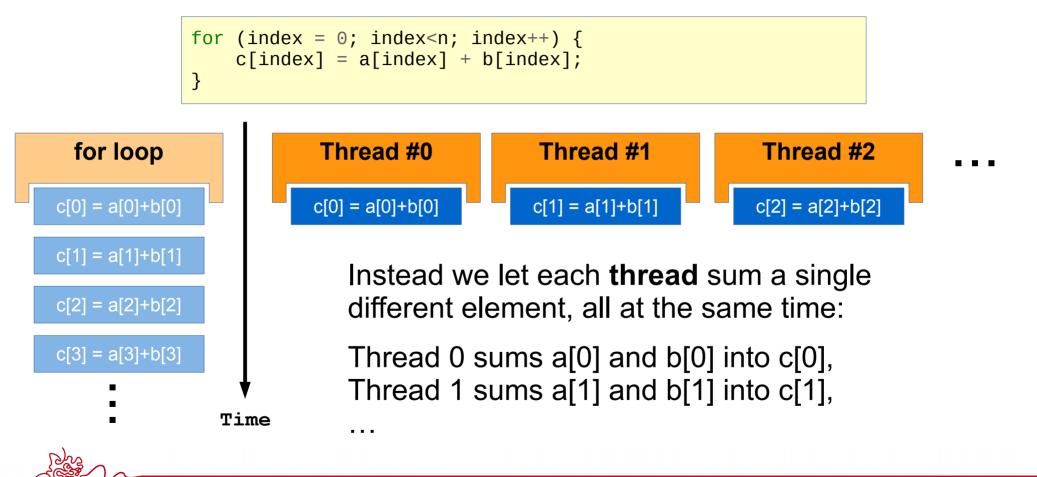
\_\_global\_\_ = can be called from host and runs on device.

Also available: \_\_device\_\_ and \_\_host\_\_.

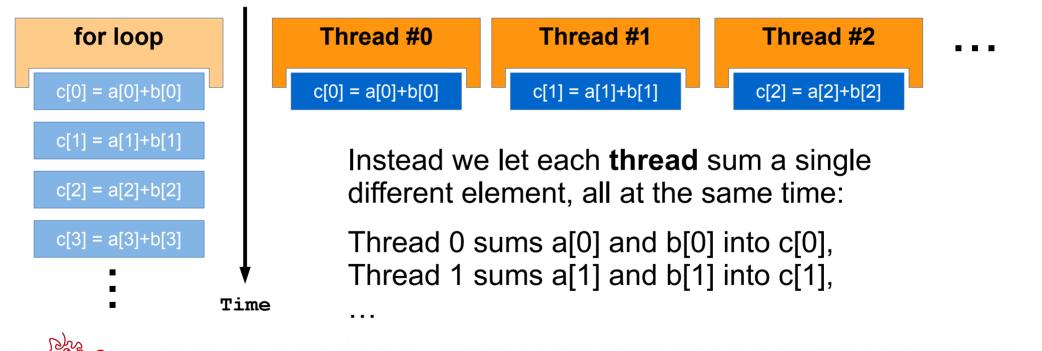
```
for (index = 0; index<n; index++) {
    c[index] = a[index] + b[index];
}</pre>
```

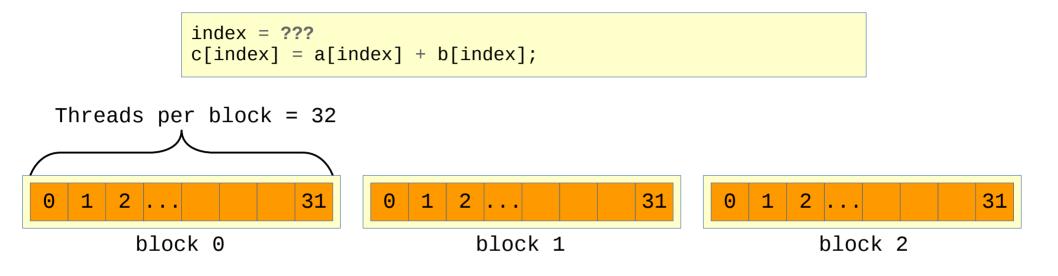


Our vector addition function steps through the loop over time, and sums a different set of elements at each time step.



index = ???
c[index] = a[index] + b[index];



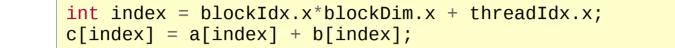


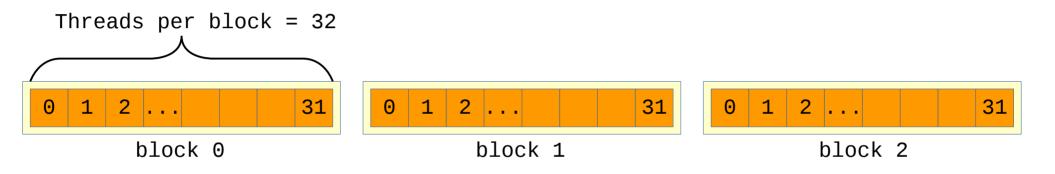
In kernel functions, CUDA automagically defines variables with the block and thread IDs.

blockDim.x = number of threads/block

blockIdx.x = current block (0, ...)

threadIdx.x = thread in current block



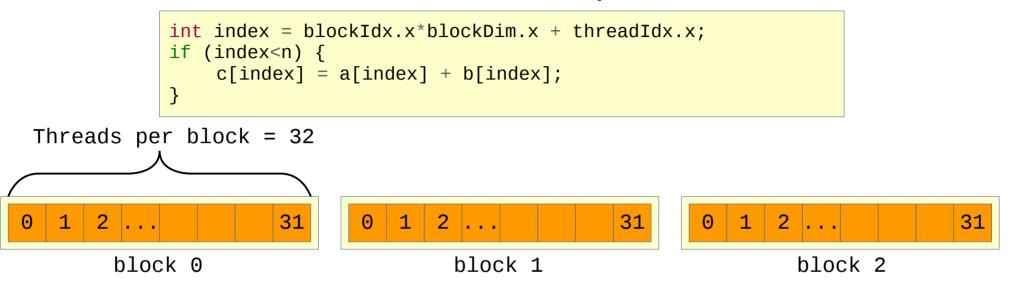


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#### Call the kernel

```
#define THREADS 512
```

```
vec_add<<< (N+THREADS-1)/THREADS, THREADS >>>(da, db, dc,N);
```

The notation is:

func<<< blocks, threads per block >>>();

"<<< , >>>" is a CUDA extension for calls to a CUDA kernel.

- allocates the blocks and threads per block that we specify;
- copy the code to the GPU card;
- set blockDim, blockIdx and threadIdx for the kernel function

#### Build and run!

- To build our CUDA program we need the compiler. Load the cuda module:
- \$ module load cuda/8.0.27
- Compile using nvcc:

```
$ nvcc -o addcuda vec_add.cu
```

• Run:

```
$ srun -p gpu --mem=1G --gres=gpu -t 1:00 ./addcuda
Time: 184.37 μs
```

- Runtime (this time): 184 µs
- Our CPU version took ~1450 µs a 7× speedup!

## But...

We're being a little unfair to the CPU. Let's see:

```
cudaMemcpy (da, a, size, cudaMemcpyHostToDevice );
cudaMemcpy (db, b, size, cudaMemcpyHostToDevice );
clock_gettime(CLOCK_REALTIME, &tic);
vec_add<<< (N+THREADS-1)/THREADS, THREADS >>>(da,db,dc,N);
cudaDeviceSynchronize();
clock_gettime(CLOCK_REALTIME, &toc);
cudaMemcpy (c, dc, size, cudaMemcpyDeviceToHost );
```

We measure the vec\_add() time, but not the memory copying:
 Only vec\_add(): ~1450 µs (CPU) ~180 µs (GPU)

## But...

Include the cudaMemcpy() calls in our time measurement:

```
clock_gettime(CLOCK_REALTIME, &tic);
cudaMemcpy (da, a, size, cudaMemcpyHostToDevice );
cudaMemcpy (db, b, size, cudaMemcpyHostToDevice );
vec_add<<< (N+THREADS-1)/THREADS, THREADS >>>(da,db,dc,N);
cudaMemcpy (c, dc, size, cudaMemcpyDeviceToHost );
clock_gettime(CLOCK_REALTIME, &toc);
```

• Measure the vec\_add() time and memory transfers:

vec\_add()+copy: ~1450 µs (CPU) ~**15000** µs (GPU) Only vec\_add(): ~1450 µs (CPU) ~**180** µs (GPU)



### Lesson:

vec\_add()+copy: ~1450 µs (CPU) ~**15000** µs (GPU) Only vec\_add(): ~1450 µs (CPU) ~**180** µs (GPU)

- GPU jobs should be compute-bound.
  - Lots of math, few memory transfers
- Memory transfers are **expensive** 
  - Do as much as possible on the GPU without transferring data
  - You can run multiple kernels after one another without moving data.



### Example #2: dot product

Dot product: p = a1\*b1 + a2\*b2 + a3\*b3 + ...

Two operations:

- Elementwise product
  - We already (almost) did in the last example
- Reduction
  - combine all elements into a single value with some function
  - Very common operation, not always trivial.



### Serial version

The serial version (dotprod.c):

- pairwise multiplication:
   tmp = a\*b
- Pairwise summation of tmp: dotp = sum(tmp)

And in main():

• We just call vec\_dot().

```
void vec_dot(float *a, float *b, float *dotp) {
    int index, s, i;
    float tmp[N];
    for (index = 0; index<N; index++) {</pre>
        tmp[index] = a[index] * b[index];
    }
    for(s = (N/2); s>0; s/=2) {
        for(i = 0; i<s; i++) {</pre>
            tmp[i] += tmp[i+s];
    *dotp = tmp[0];
```

float dotp = 0.0; vec\_dot(a, b, &dotp);

### **Pairwise Summation**

Summing large data sets naively will cause a form of *catastrophic cancellation* — one term will become orders of magnitude larger than the other, and you lose significant digits.

```
for(s = (N/2); s>0; s/=2) {
    for(i = 0; i<s; i++) {
        tmp[i] += tmp[i+s];
     }
}</pre>
```

We sum elements in pairs:

```
tmp[0] += tmp[N/2]; tmp[1] += tmp[N/2+1]...
```

Then sum the pairs in pairs until we have a single element. All terms now have the same order of magnitude.

## **Pairwise Summation**

Summing large data sets naively will cause catastrophic cancellation

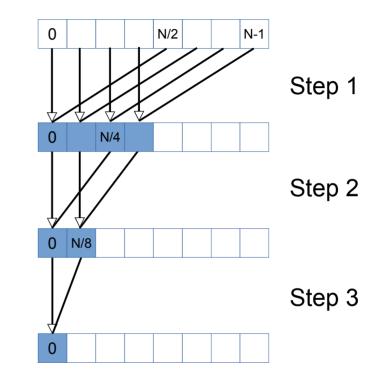
We avoid it by pairwise summation:

We sum elements in pairs:

tmp[0] = tmp[0] + tmp[N/2];tmp[1] = tmp[1] + tmp[N/2+1]...

We recursively sum each pair in the same way, until we have a single element.

All summation terms will have about the same order of magnitude.





## Run the serial code

Compile our example program:

\$ gcc -03 dotprod.c -o dotpcpu

Run it on Tombo:

\$ srun -t 1:00 --mem=50m ./dotpcpu
calculated dot product: 26214400.00
time: 4899.19 μs

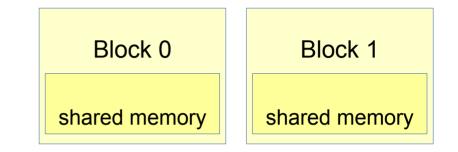
We get times in the 4800-5200µs range



Blocks have private "shared" memory

- 48KB (can be changed)
- faster access than global memory
- memory accesses in different blocks are independent
  - $\rightarrow$  no memory contention

- Read input from global memory
- Calculate, using shared memory for intermediate values
- each block does a partial reduction one partial value per block
- finally add its partial value to the final result in global memory

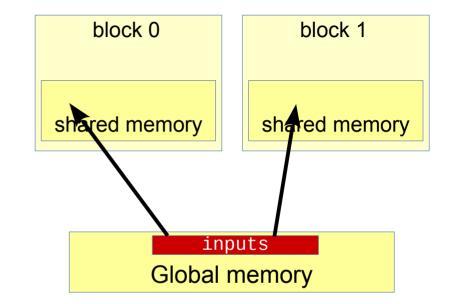




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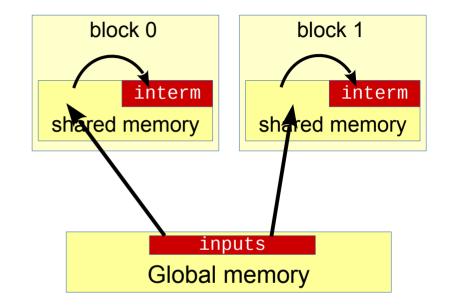
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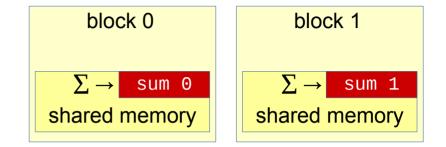
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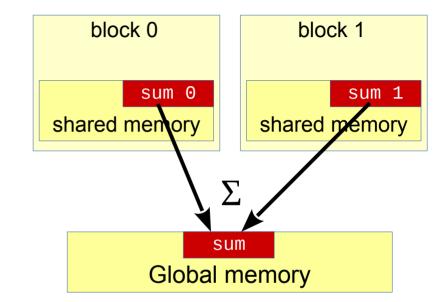




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The GPU version (dotprod.cu), elementwise multiplication:

```
__global__ void vec_dotp(float *a, float *b, float *dotp) {
    int s;
    int tid = threadIdx.x;
    int index = blockIdx.x*blockDim.x + threadIdx.x;
    // allocate block-local memory
    __shared__ float tmp[THREADS];
    tmp[tid] = a[index] * b[index];
    __syncthreads();
```

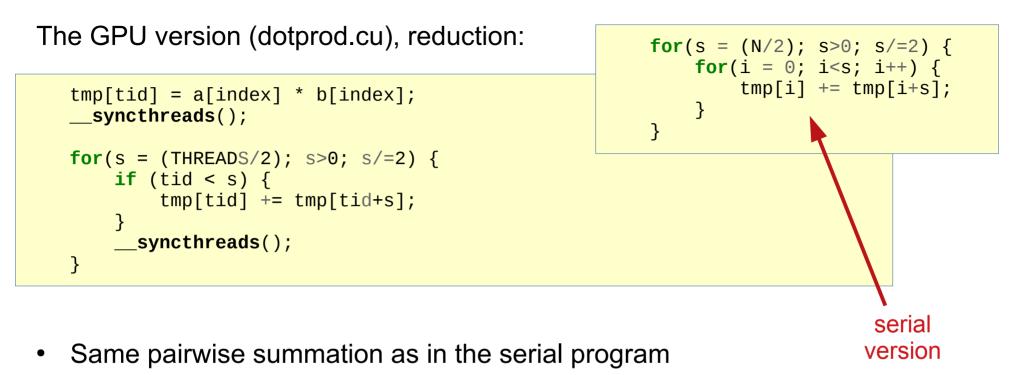
- Get local memory with "\_\_\_shared\_\_\_".
  - Much faster than global memory, but limited size
- \_\_\_\_\_syncthreads() synchronizes all threads.
  - Threads in a warp are synchronized, but threads in *different* warps are not.

The GPU version (dotprod.cu), reduction:

```
tmp[tid] = a[index] * b[index];
__syncthreads();
for(s = (THREADS/2); s>0; s/=2) {
    if (tid < s) {
        tmp[tid] += tmp[tid+s];
    }
    __syncthreads();
}
```

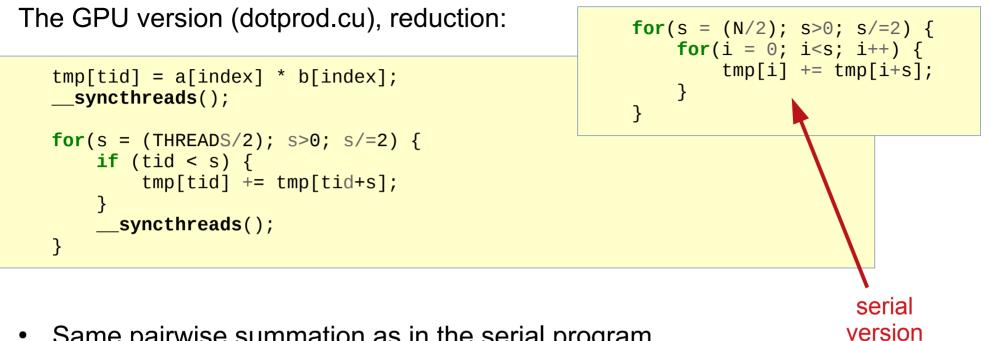
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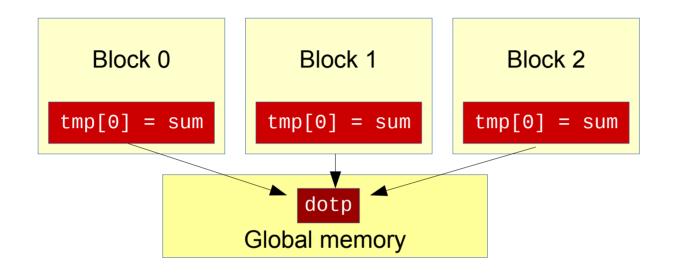




- Same pairwise summation as in the serial program
  - But inner loop is parallel
  - We must synchronize after each iteration so all threads really are finished.
- **NOTE:** this example works only when THREADS is a power of 2 and N is a multiple of THREADS



We have the partial results in each block:



- Add the partial sums together into the final result.
- But blocks are independent, and could access dotp at the same time (*"race condition"*)

 $\rightarrow$  need to use an "atomic" operation: atomicAdd()

The complete reduction:

```
for(s = (THREADS/2); s>0; s/=2) {
    if (tid < s) {
        tmp[tid] += tmp[tid+s];
    }
    ___syncthreads();
}
if (tid == 0) {
    atomicAdd(dotp, tmp[0]);
}</pre>
```

- Pairwise sum all elements in the block
- Finally thread #0 atomically adds the result into the return parameter



• Compile:

\$ nvcc -o dotpcuda dotprod.cu

• Run it on Tombo:

\$ srun -p gpu --mem=50m --gres=gpu -t 1:00 ./dotpcuda calculated dot product: 26214400.000000 time: 5258.94 μs

Result: ~4800-5200 µs (CPU) ~5100-5500 µs (GPU)

#### Let's try adding more work: run the kernel twice

CPU call:

// calculate dot product of a and b, return in dotp
vec\_dot(a, b, &dotp);
vec\_dot(a, b, &dotp);

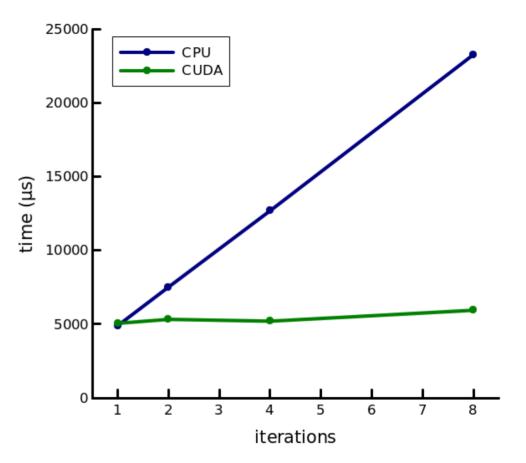
#### CUDA call:

// calculate dot product of a and b, return in dotp
vec\_dot<<< (N+THREADS-1)/THREADS, THREADS>>>(da, db, ddotp);
vec\_dot<<< (N+THREADS-1)/THREADS, THREADS>>>(da, db, ddotp);

```
Result: ~7500 µs (CPU) ~5300 µs (GPU)
```

## **CPU and CUDA**

- Let's test for different number of iterations of our dot product:
- CUDA has large, fixed transfer cost
   → For small amounts of work, a single CPU core is faster
  - Even for 8 dot products, the time is dominated by the data transfer
- Tombo nodes have 16 cores. If we use them, the CPU will be faster.
- Practical GPU speed improvement is usually less than 3-4 times CPU.



## Summary

- Use \_\_shared\_\_ to allocate shared memory
  - Fast, but limited size (~48Kb)
  - Does not persist between different kernels
- \_\_\_\_\_syncthreads() synchronizes all threads in a block
  - threads in a warp are synchronous, but different warps are not.
- Use atomic operations when multiple threads have to change the same data
  - Blocks are independent, so atomics are necessary
  - Atomic operations in global memory faster than in shared.



## **Final Points**

- GPU computation is fast. GPU data transfer is slow.
  - To reduce transfer amount, filter data on the CPU.
  - Do as much calculation as possible on the data in the GPU.
  - Avoid storing intermediate values on the host.
  - Running kernels is cheap
- Memory organisation matters a lot
  - keep data in block shared memory or thread local memory
  - Access data sequentially
  - Global memory is persistent across kernels



## Odds and Ends

You can write functions that run on both host and device:

```
__HOST__ __DEVICE__ float myfunc() {
#ifdef __CUDA_ARCH__
    // CUDA code, probably called from a kernel
#else
    // Host code, running without CUDA
#endif
}
```

\_\_CUDA\_ARCH\_\_ defines the compute capability level, but is only defined in code that runs on the GPU.



# Odds and Ends

• NVIDIA separates it's cards by "compute capability". Each newer capability version is a superset of previous ones.

Set the desired capability level with "--arch" parameter:

\$ nvcc --arch=sm\_35 -o dotpcuda dotprod.cu

- For example: atomicAdd() for integers appeared in 1.1 (sm\_11); floating point version in 2.0 (sm\_20) and double precision only in version 6.0 (sm\_60).
- Wikipedia has a great page: https://en.wikipedia.org/wiki/CUDA

