# An Introduction to CUDA and GPU Programming

Scientific Computing and Data Analysis



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# GPU

- GPU is a "Graphics Processing Unit"
	- Developed for games
	- High computational capability



- $\cdot$  CUDA is NVIDIAs C/C++ language extension for HPC computing on GPUs
	- But other options exist: OpenCL, OpenACC...

# Trend: multiple cores, parallel execution

As transistor count grows, CPUs have gained more cores and more features.

**A CPU** is a collection of 2-20 independent, fully **general cores** connected in a local network, with local and shared memory, and I/O to the outside world.

Xeon E5 2699v3

5.5B transistors 18 cores 45 MB L3 cache



# Meanwhile...

- PC graphics gained hardware support, 2D then 3D
- Early 2000s: GPUs with on-card programmable 3D transformation and lighting calculations ("shaders")

A shader is a *small* program that:

- Runs on many pixels or many vertices at once - ("SIMD")
- Is compute bound only math with no I/O or conditionals
- Is stream oriented can process a stream of data points with little or no additional state



# Meanwhile...

- Researchers realized that you can run any function as a shader First only a curiosity, but it became obvious that GPUs were potentially powerful general code accelerators.
- 2007: NVIDIA releases **CUDA**, Apple releases **OpenCL** (later standardized under Chronos).

Both are direct means of harnessing the GPU compute power without hacks and workarounds.





#### CPU programs are **complex**:

• OS, simulators, games...

#### Cores are **independent**:

• Each core is a full computer

#### **General** computing tasks:

• Math, I/O, interactive applications.

#### GPU programs are **simple**:

• Small size, simple control flow

#### Cores are very **parallel**:

• Many execution threads share one control unit

#### Very **specialized** for math:

• floating point function evaluators.



- The CPU ("*host*") and the GPU ("*device*") are separate  $\rightarrow$  you need to copy data to and from the GPU
- The GPU has "streaming multiprocessors" = processors that each can run hundreds of threads.
	- Each SM has 8-64 CUDA cores; with more threads they take turns on the SM

# GPU concepts

A **thread** runs a single computation.

- Like CPU threads, they share memory and code with other threads.
- Much simpler, slower than CPU cores.
- Limited thread-local memory, registers.



### A **warp** is a collection of 32 threads

– All threads in a warp run the *same code* at the *same time*  $\rightarrow$  one thread takes less space and energy than a CPU core

- Threads have their own registers and variables.
- Same instructions, but act on different data ("SIMD")



# Blocks

- A **block** is a collection of up to 1024 threads
	- All threads run the same code
	- Has shared fast memory (48KB)
- You can organize threads in 1-D, 2-D or 3-D, but that's only programmer convenience



# Grid

#### A **grid** is a set of blocks

- Blocks are **independent**
	- Can not access data in shared memory of other blocks
- Blocks run in **any order.**
- Threads in all blocks have main memory in common
- You can organize blocks in different dimensions, the same as threads in a block.







# Workflow

- 1) Write a C function (a "*kernel*") that will run on all **threads**
- 2) Copy your data from **host** to **device**
- 3) Run the **kernel** on the **device** with your data, using some number of **blocks** and **threads**
- 4) Copy results from **device** back to **host**
- 5) Repeat from 3); or finish up
- 6) Done!



# Log In on Tombo!

• First, let's all log in on Tombo:

\$ ssh <your-ID>@tombo.oist.jp

• Copy the code from the common area:

\$ cp -r /work/training/GPU .

• Go to the new directory:

#### \$ cd GPU/code

## GPU Resources at OIST

#### **Tombo:**

"gpu" - 1 node 2\* K40 Training

#### **Sango:**

"gpu" - 2 node 4\* K80 image analysis pipelines "powernv" - 4 node 4\* P100 general, deep learning

#### **Saion:**

"gpu" - 9 node 4\* P100 general, deep learning - 8 node 4\* V100 "powernv" - 4 node 4\* V100 general, deep learning

\$ ssh <your-ID>@tombo.oist.jp \$ cp -r /work/share/training/GPU . \$ cd GPU/code

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• You need to ask for the GPU partition, and reserve the GPU resource:

\$ srun --partition=gpu --gres=gpu <program>

- Use multiple cards with subsest - gres=gpu: N
- For compilation, load the cuda module:

\$ module load cuda/8.0.27

### Program #1: GPU vector addition

We will adapt a simple program to run on the GPU:

```
void vec_add(float *a, float *b, float *c, int n) {
     int index;
   for (index=0; index<n; index++) {
        c[index] = a[index] + b[index];
     }
}
```
Vector addition is the "hello world" of parallel programming The source is GPU/code/vec\_add.c

The CPU-based version

- Allocate three arrays:
	- a, b inputs c output
- Initialize all elems
	- $a, b = 5.0$  $c = 0.0$
- Call vec\_add()

```
void vec_add(float *a, float *b, float *c, int n) {
     int index;
    for (index=0; index<n; index++) {
        c[index] = a[index] + b[index];
     }
}
int main(int argc, char **argv) {
    int i;
    float *C, *b, *a;c = (float *)<b>malloc</b>( size);b = (float *)<b>malloc</b>( size);a = (float *)<b>malloc</b>( size);for (i = \theta; i<N; i++ ) {
        a[i] = b[i] = 5.0;c[i] = 0.0; }
     vec_add(a,b,c,N);
}
```
#### We have two extra things in our code. Timing:

```
double time_diff_nano(struct timespec *toc, struct timespec *tic) {
     return (1e9*(toc->tv_sec-tic->tv_sec)+
         (toc->tv_nsec-tic->tv_nsec));
}
...
clock_gettime(CLOCK_REALTIME, &tic);
...
clock_gettime(CLOCK_REALTIME, &toc);
time_diff_nano(&toc, &tic)/1000.0;
```
#### and sanity check:

```
float sum = 0.0;
for (index = \theta; index<N; index++) {
    sum += c[index];}
if (fabs((sum-CORRECT)/sum)>EPSILON) {
     printf(" correct sum: %.2f - fail\n", CORRECT);
}
```
### Run the serial code

• Compile our example program:



\$ srun -t 1:00 --mem=50m ./addcpu calculated sum: 10485760.00 Time: 1441.77 µs

• Time: ~1425-1550µs

### Allocate memory on the GPU

- Make a copy (or just use code/vec add.cu) and open it: \$ cp vec add.c vec add.cu
- Allocate memory on the GPU (and free it at the end):

```
float *dc, *db, *da;
cudaMalloc( &da, size );
cudaMalloc( &db, size );
cudaMalloc( &dc, size );
cudaFree( da );
cudaFree( db );
cudaFree( dc );
```
- All cuda functions start with "cuda"
- $-$  You usually don't need to add  $\#inc$  lude statements the cuda compiler adds then for you.

#### Move our data to the GPU

Copy our data over:



cudaMemcpyHostToDevice = copy to GPU from computer cudaMemcpyDeviceToHost = copy to computer from GPU

also cudaMemcpyHostToHost,cudaMemcpyDeviceToDevice

Many function variations available: **cudaMemcpyAsync**(), **cudaMemcpy2DTo**[**From**]**Array**(), **cudaMemcpyToSymbol**() …

#### Create the kernel

Our original vector addition function:

```
void vec_add(double *a, double *b, double *c, int n) {
     int index;
    for (index = \theta; index<n; index++) {
        cfindex] = afindex] + bfindex];
 }
}
```
#### Add "\_ \_global\_ \_" specifier:

 $\Box$ global $\Box$  void vec\_add(double \*a, double \*b, double \*c, int n) {

\_ \_global\_ \_ = can be called from **host** and runs on **device**.

Also available:  $\_\_$  device  $\_\_$  and  $\_\_$  host  $\_\_$ .

```
for (index = 0; index<n; index++) {
    cfindex] = a[index] + b[index];
}
```


Our vector addition function steps through the loop over time, and sums a different set of elements at each time step.



index = **???**  $c$ [index] = a[index] + b[index];





In kernel functions, CUDA automagically defines variables with the block and thread IDs.

 $b$ lockDim. $x =$  number of threads/block

 $blockIdx.x = current block (0, ...)$ 

 $threadIdx.x = thread in current block$ 

int index =  $blockIdx.x*blockDim.x + threadIdx.x;$  $c$ [index] = a[index] + b[index];



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#### Call the kernel

```
#define THREADS 512
```

```
vec_add<<< (N+THREADS-1)/THREADS, THREADS >>>(da, db, dc,N);
```
The notation is:

func<<< blocks, threads per block >>>();

"<<< , >>>" is a CUDA extension for calls to a CUDA kernel.

- allocates the blocks and threads per block that we specify;
- copy the code to the GPU card;
- set blockDim, blockIdx and threadIdx for the kernel function

#### Build and run!

- To build our CUDA program we need the compiler. Load the cuda module:
- \$ module load cuda/8.0.27
- Compile using nvcc:
	- \$ nvcc -o addcuda vec add.cu
- Run:

```
$ srun -p gpu --mem=1G --gres=gpu -t 1:00 ./addcuda 
 Time: 184.37 µs
```
- Runtime (this time): 184 µs
- Our CPU version took ~1450 µs a **7×** speedup!

# **But**

#### We're being a little unfair to the CPU. Let's see:

```
 cudaMemcpy (da, a, size, cudaMemcpyHostToDevice );
 cudaMemcpy (db, b, size, cudaMemcpyHostToDevice );
 clock_gettime(CLOCK_REALTIME, &tic);
vec_add<<< (N+THREADS-1)/THREADS, THREADS >>>(da,db,dc,N);
 cudaDeviceSynchronize();
 clock_gettime(CLOCK_REALTIME, &toc);
 cudaMemcpy (c, dc, size, cudaMemcpyDeviceToHost );
```
• We measure the vec\_add() time, but not the memory copying: Only vec\_add(): ~1450 µs (CPU) ~180 µs (GPU)

# **But**

Include the cudaMemcpy() calls in our time measurement:

```
 clock_gettime(CLOCK_REALTIME, &tic);
 cudaMemcpy (da, a, size, cudaMemcpyHostToDevice );
 cudaMemcpy (db, b, size, cudaMemcpyHostToDevice );
vec_add<<< (N+THREADS-1)/THREADS, THREADS >>>(da,db,dc,N);
 cudaMemcpy (c, dc, size, cudaMemcpyDeviceToHost );
 clock_gettime(CLOCK_REALTIME, &toc);
```
• Measure the vec\_add() time *and* memory transfers:

Only vec\_add(): ~1450 µs (CPU) ~**180** µs (GPU) vec\_add()+copy: ~1450 µs (CPU) ~**15000** µs (GPU)

### Lesson:

Only vec\_add(): ~1450 µs (CPU) ~**180** µs (GPU) vec\_add()+copy: ~1450 µs (CPU) ~**15000** µs (GPU)

- GPU jobs *should be* compute-bound.
	- Lots of math, few memory transfers
- Memory transfers are **expensive**
	- Do as much as possible on the GPU without transferring data
	- You can run multiple kernels after one another without moving data.

### Example #2: dot product

Dot product:  $p = a1 * b1 + a2 * b2 + a3 * b3 + ...$ 

Two operations:

- Elementwise product
	- We already (almost) did in the last example
- *Reduction*
	- combine all elements into a single value with some function
	- Very common operation, not always trivial.

### Serial version

The serial version (dotprod.c):

- pairwise multiplication:  $tmp = a * b$
- Pairwise summation of tmp:  $dot p = sum(tmp)$

And in main():

• We just call vec\_dot().

```
void vec_dot(float *a, float *b, float *dotp) {
    int index, s, i;
     float tmp[N];
    for (index = \theta; index<N; index++) {
        tmp[index] = a[index] * b[index]; }
    for(s = (N/2); s>0; s/=2) {
        for(i = 0; i<s; i++) {
            tmp[i] += tmp[i+s]; }
 }
    \stardotp = tmp[0];
}
```
float  $dot p = 0.0$ ; vec\_dot(a, b, &dotp);

### Pairwise Summation

Summing large data sets naively will cause a form of *catastrophic cancellation* — one term will become orders of magnitude larger than the other, and you lose significant digits.

```
for(s = (N/2); s>0; s/=2) {
        for(i = 0; i < s; i + 1) {
            tmp[i] += tmp[i+s]; }
 }
```
We sum elements in pairs:

 $tmp[0] += tmp[N/2]; tmp[1] += tmp[N/2+1]...$ 

Then sum the pairs in pairs until we have a single element. All terms now have the same order of magnitude.

### Pairwise Summation

Summing large data sets naively will cause *catastrophic cancellation*

We avoid it by pairwise summation:

We sum elements in pairs:

 $tmp[0] = tmp[0] + tmp[N/2];$  $tmp[1] = tmp[1] + tmp[N/2+1]...$ 

We recursively sum each pair in the same way, until we have a single element.

All summation terms will have about the same order of magnitude.



### Run the serial code

Compile our example program:

\$ gcc -O3 dotprod.c -o dotpcpu

Run it on Tombo:

\$ srun -t 1:00 --mem=50m ./dotpcpu calculated dot product: 26214400.00 time: 4899.19 µs

We get times in the 4800-5200 µs range

Blocks have private "shared" memory

- 48KB (can be changed)
- faster access than global memory
- memory accesses in different blocks are independent
	- $\rightarrow$  no memory contention

- Read input from global memory
- Calculate, using shared memory for intermediate values
- each block does a partial reduction one partial value per block
- finally add its partial value to the final result in global memory





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The GPU version (dotprod.cu), elementwise multiplication:

```
__global__ void vec_dotp(float *a, float *b, float *dotp) {
   int s;
  int tid = threadIdx.x;
  int index = blockIdx.x*blockDim.x + threadIdx.x; // allocate block-local memory
   __shared__ float tmp[THREADS];
  tmp[tid] = a[index] * b[index]; __syncthreads();
```
- Get local memory with "\_\_shared\_\_".
	- Much faster than global memory, but limited size
- $\_\_$ syncthreads() synchronizes all threads.
	- Threads in a warp are synchronized, but threads in *different* warps are not.

The GPU version (dotprod.cu), reduction:

```
tmp[tid] = a/index] * b[index]; __syncthreads();
   for(s = (THEADS/2); s>0; s/=2) {
       if (tid < s) {
           tmp[tid] += tmp[tid+s]; }
        __syncthreads();
 }
```
- Same pairwise summation as in the serial program
	- But inner loop is parallel
	- We must synchronize after each iteration so all threads really are finished.



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- Same pairwise summation as in the serial program
	- But inner loop is parallel
	- We must synchronize after each iteration so all threads really are finished.
- **NOTE**: this example works only when THREADS is a power of 2 and N is a multiple of THREADS

We have the partial results in each block:



- Add the partial sums together into the final result.
- But blocks are independent, and could access dotp at the same time (*"race condition"*)

→ need to use an "*atomic*" operation: atomicAdd()

#### The complete reduction:

```
for(s = (THREADS/2); s>0; s/=2) {
       if (tid \leq s) {
           tmp[tid] += tmp[tid+s]; }
        __syncthreads();
    }
   if (tid == 0) {
        atomicAdd(dotp, tmp[0]);
 }
```
- Pairwise sum all elements in the block
- Finally thread #0 atomically adds the result into the return parameter

• Compile:

\$ nvcc -o dotpcuda dotprod.cu

• Run it on Tombo:

\$ srun -p gpu --mem=50m --gres=gpu -t 1:00 ./dotpcuda calculated dot product: 26214400.000000 time: 5258.94 µs

Result: ~4800-5200 µs (CPU) **~5100-5500 µs** (GPU)

#### Let's try adding more work: run the kernel twice

CPU call:

 // calculate dot product of a and b, return in dotp vec\_dot(a, b, &dotp); vec\_dot(a, b, &dotp);

#### CUDA call:

 // calculate dot product of a and b, return in dotp vec\_dot<<< (N+THREADS-1)/THREADS, THREADS>>>(da, db, ddotp); vec dot<<< (N+THREADS-1)/THREADS, THREADS>>>(da, db, ddotp);

Result: ~7500 µs (CPU) **~5300 µs** (GPU)

### CPU and CUDA

- Let's test for different number of iterations of our dot product:
- CUDA has large, fixed transfer cost
	- $\rightarrow$  For small amounts of work, a single CPU core is faster
	- Even for 8 dot products, the time is dominated by the data transfer
- $\cdot$  Tombo nodes have 16 cores. If we use them, the CPU will be faster.
- Practical GPU speed improvement is usually less than 3-4 times CPU.



## Summary

- Use \_\_shared \_\_ to allocate shared memory
	- Fast, but limited size (~48Kb)
	- Does not persist between different kernels
- $\_\_$  syncthreads () synchronizes all threads in a block
	- threads in a warp are synchronous, but different warps are not.
- Use atomic operations when multiple threads have to change the same data
	- Blocks are independent, so atomics are necessary
	- Atomic operations in global memory faster than in shared.

### Final Points

- GPU *computation* is fast. GPU *data transfer* is slow.
	- To reduce transfer amount, filter data on the CPU.
	- Do as much calculation as possible on the data in the GPU.
	- Avoid storing intermediate values on the host.
	- Running kernels is cheap
- Memory organisation matters a lot
	- keep data in block shared memory or thread local memory
	- Access data sequentially
	- Global memory is persistent across kernels

### Odds and Ends

You can write functions that run on both host and device:

```
_ _HOST_ _ _ _DEVICE__ float myfunc() {
\# \texttt{ifdef} CUDA ARCH
    // CUDA code, probably called from a kernel
#else
    // Host code, running without CUDA
#endif
}
```
 $\mathsf{L}_\mathsf{L}$  CUDA\_ARCH\_  $\mathsf{L}_\mathsf{L}$  defines the compute capability level, but is only defined in code that runs on the GPU.

# Odds and Ends

• NVIDIA separates it's cards by "compute capability". Each newer capability version is a superset of previous ones.

Set the desired capability level with "--arch" parameter:

\$ nvcc **--arch=sm\_35** -o dotpcuda dotprod.cu

- For example: atomicAdd() for integers appeared in 1.1 (sm\_11); floating point version in 2.0 (sm\_20) and double precision only in version 6.0 (sm\_60).
- Wikipedia has a great page: https://en.wikipedia.org/wiki/CUDA