

Computation of decomposition numbers of G_n

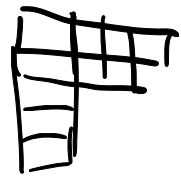
$$H_n = H_{\mathbb{F}, \mathbb{F}}(G_n), \quad \zeta \in \mathbb{F}^*, \quad \zeta = 1$$

$e = \text{mult char of } \mathbb{F}$.

Main open problem: D^M simple modules

μ e -reg partition of n ,
 \hookrightarrow no e equal parts

Ex: $\lambda = (4, 3, 1, 1, 1) \vdash 10$

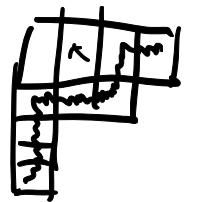


$$[S^\lambda : D^\mu] \geq 0$$

\hookrightarrow Specht mod

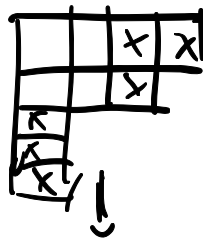
① $[S^\lambda : D^\mu]_{\mathbb{F}}^e > 0$ only if λ, μ belong to same block

- Def: - rim $(i, j) \in \lambda$ s.t. $(i+1, j+1) \notin \lambda$



- rim e -hooks:

$$e=3$$



- e -weight:



3-weight of $(4, 3, 1, 1, 1)$ is 3

= e-core $p \vdash n$ -we \Rightarrow 3-core of $(4, 3, 1, 1)$
is $\rho = (1)$

Thm: (Nakayama conj): $\lambda, \mu \vdash n$ belong to the same
(Brouer-Robinson thm)

Block \Leftrightarrow they have same e-core

- We can study decomp numbers by block (by e-core)

- e-weight measures difficulty of rep theory of the block.

- Decomp number problem solved for $w \leq 2$

• $w = 0$

• $w = 1$

• $w = 2$ Richards $p = \text{char}(\mathbb{F}) \geq 2$, Fayers, $p = 2$.

• $w \geq 3$

② LLT algorithm when $\mathbb{F} = \mathbb{C}$

$$D = \left([s^\lambda: D^\mu] \right)_{(\lambda, \mu)}, \quad D_{\mathbb{F}} = D_{\mathbb{C}} \begin{matrix} A \\ \downarrow \\ \text{Adjustment} \\ \text{matrix} \end{matrix}$$

James's Conj: $\exists \lambda \in \text{Cher}(\mathbb{F}) \Rightarrow A = \underline{I}$

(G. Williamson '17 found a counterexample)

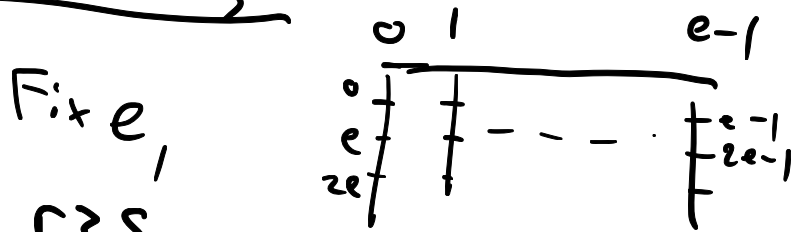
James's conj holds for $w \leq 4$, principal block of \mathcal{H}_{se} (Law).

③ $\mu, \lambda \vdash n, \mu \geq \lambda, \forall i \sum_{j=0}^i \mu_j \geq \sum_{j=0}^i \lambda_j$

$\lambda = (4, 3, 1, 1, 1), \mu = (3, 3, 2, 2)$

thm: $[S^\lambda; D^\mu] = 0$ unless $\mu \geq \lambda$

the abacus

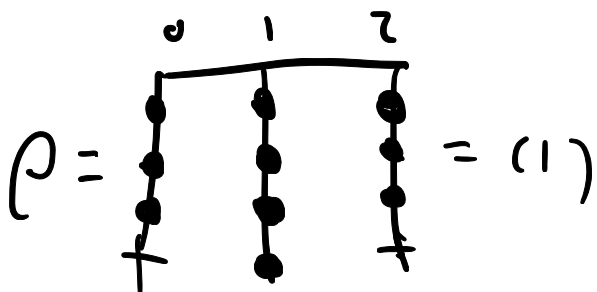
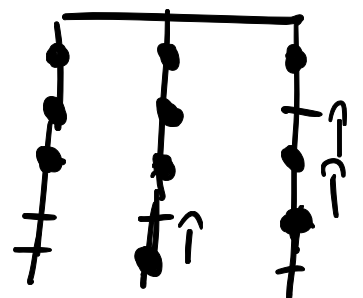


$\lambda = (\lambda_1, \dots, \lambda_s)$

$\left\{ \begin{array}{l} r-s \text{ beads in the } r-s \\ \text{first positions} \\ \text{like } \lambda_i, \text{ empty spaces} \\ \text{between each bead} \end{array} \right.$

Ex: $\lambda = (4, 3, 1, 1, 1), e = 3, r = 10$

$r-s = 5$



④ Block B
 runner i there are k more beads than
 in runner $i-1$ either;

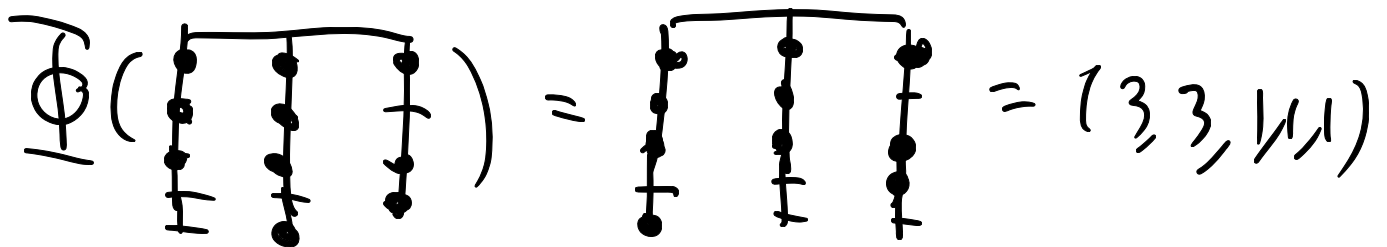
- If $k \geq w \Rightarrow \Phi: B \rightarrow \hat{B}$ s.t.

$$[S^\lambda; D^\mu] = [S^{\Phi(\lambda)}; D^{\Phi(\mu)}]$$

$\Phi(\lambda)$ build from λ by swapping runners
 $i, i-1$.

- If $k < w \Rightarrow$ If in runner i in λ there are
 more than k beads with empty space to their left,
 then λ exceptional

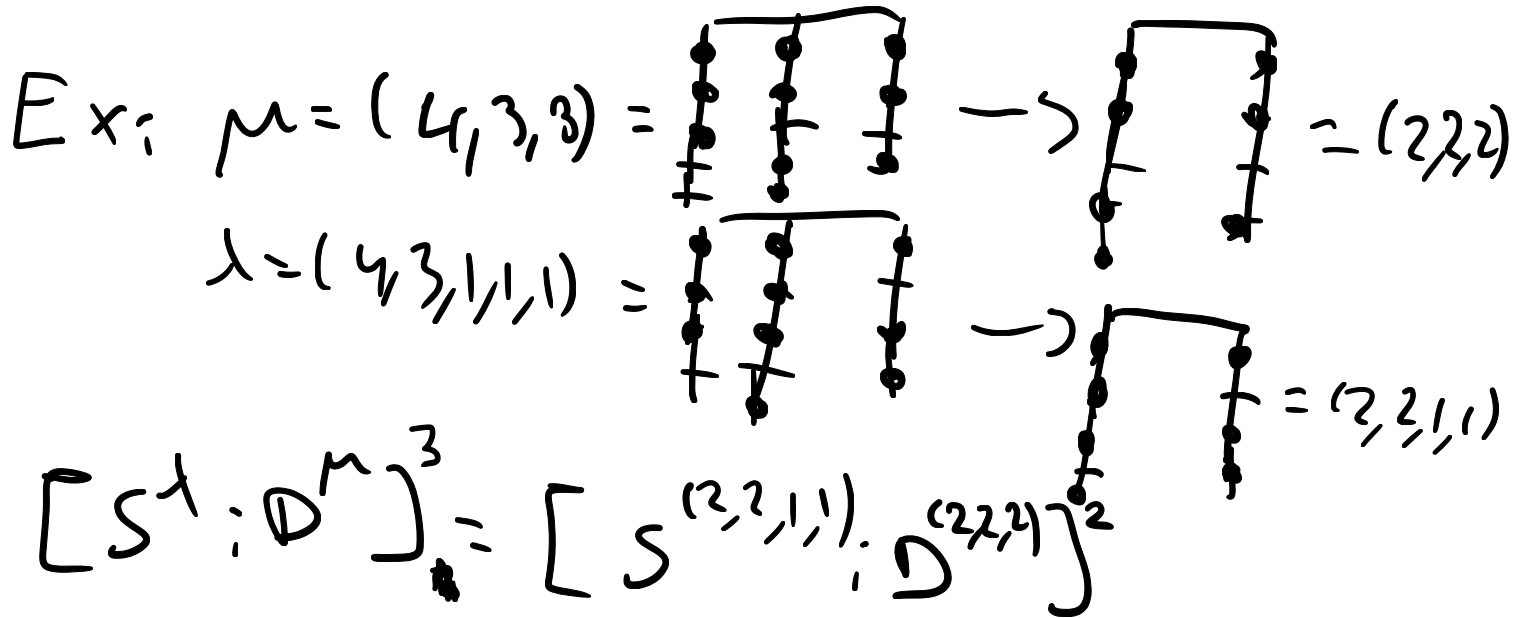
ex: $[S^{(4,3,1,1)}; D^\mu] = [S^{(3,3,1,1)}; D^{\Phi(\mu)}]$



⑤ $\lambda, \mu \in B$, runner i define a function

$L_i(\lambda) = L_i(\mu)$ and λ, μ i -empty

$$\Rightarrow [S^\lambda; \mathbb{D}^\mu]^e = [S^{\lambda^i}; \mathbb{D}^{\mu^i}]^{e-1}$$



- M. Rees's max comp length principal block of VF_{3p}^1 is 14, ($w=3$)