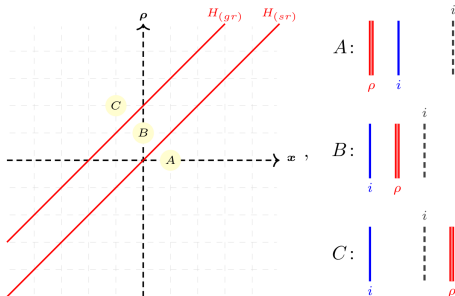


# On weighted KLRW algebras

Or: Mind the distance

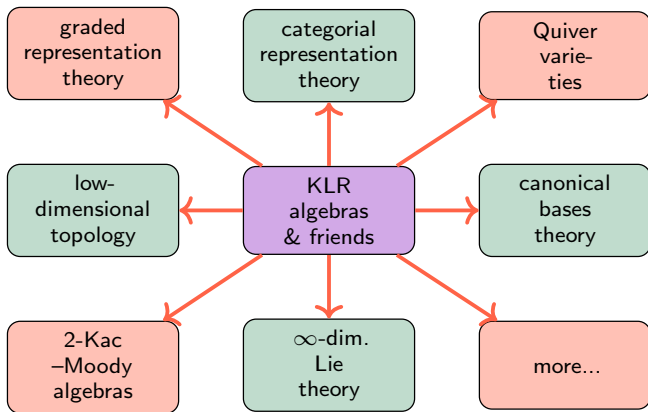
Daniel Tubbenhauer



Joint with Andrew Mathas

February 2022

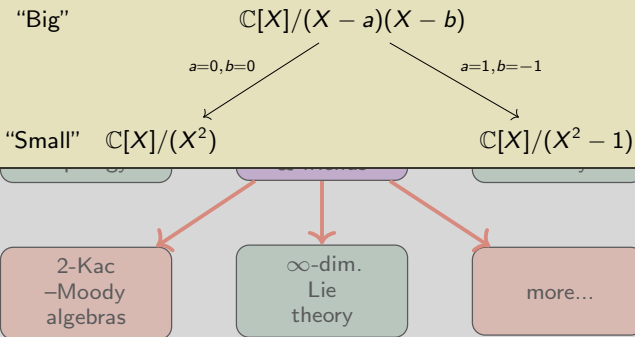
## Where are we?



- ▶ **Khovanov–Lauda–Rouquier ~2008 + many others (including OIST)**  
KLR algebras are at the heart of categorical representation theory
- ▶ Similarly for quiver Schur algebras and diagrammatic Cherednik algebras
- ▶ **Problem** All of these are actually really complicated!

## Observation

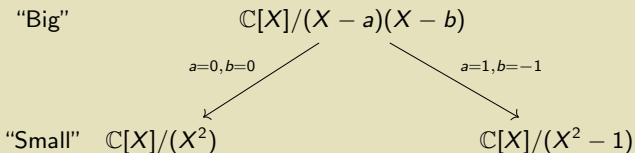
It often helps to find a “bigger” interpolating algebra, e.g.:



- ▶ **Khovanov–Lauda–Rouquier ~2008 + many others (including OIST)**  
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- ▶ Similarly for quiver Schur algebras and diagrammatic Cherednik algebras
- ▶ **Problem** All of these are actually really complicated!

## Observation

It often helps to find a “bigger” interpolating algebra, e.g.:



## Today

How to play the interpolation game using **planar geometry**?

As an upshot we get an algebra interpolating between various algebras appearing in categorical representation theory

The takeaway keyword: **Distance!**

- ▶ Kh
- ▶ KL
- ▶ Sim

(ST)

ras

- ▶ **Problem** All of these are actually really complicated!

## String diagrams – the baby case

---

Connect eight points at the bottom with eight points at the top:



or

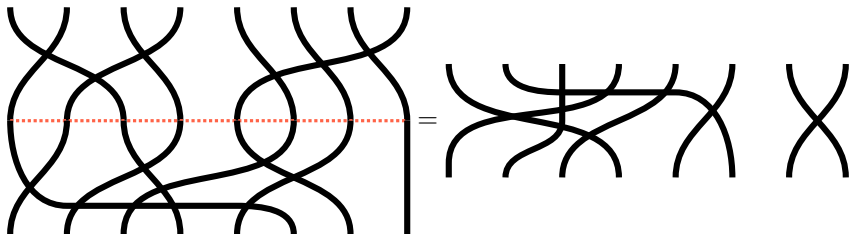


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We just invented the symmetric group  $S_8$

## String diagrams – the baby case

---

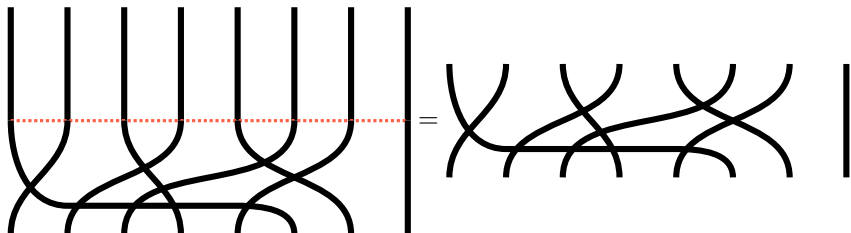


---

My multiplication rule for  $gh$  is “stack  $g$  on top of  $h$ ”

## String diagrams – the baby case

- ▶ We clearly have  $g(hf) = (gh)f$
- ▶ There is a do nothing operation  $1g = g = g1$



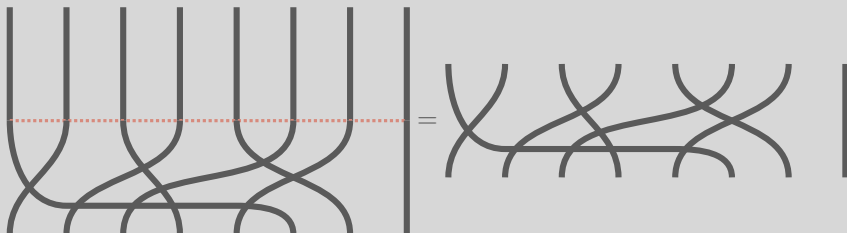
- ▶ Generators–relations (the Reidemeister moves)



The bait

In diagram algebras relations, properties, etc. become visually clear

- ▶ We clearly have
- ▶ There is a do nothing operation  $1g = g = g1$



- ▶ Generators–relations (the Reidemeister moves)





## The bait

In diagram algebras relations, properties, etc. become visually clear

- ▶ We clearly have
- ▶ There is a do nothing operation  $1g = g = g1$

## The catch

Diagram algebras are usually “not really” using any planar geometry

For example, the diagrams for symmetric groups are just algebra written differently

- ▶ Generators–relations (the Reidemeister moves)



### The bait

In diagram algebras relations, properties, etc. become visually clear

- ▶ We clearly have
- ▶ There is a do nothing operation  $1g = g = g1$

### The catch

Diagram algebras are usually “not really” using any planar geometry

For example, the diagrams for symmetric groups are just algebra written differently

- ▶ General

### Idea (Webster ~2012)

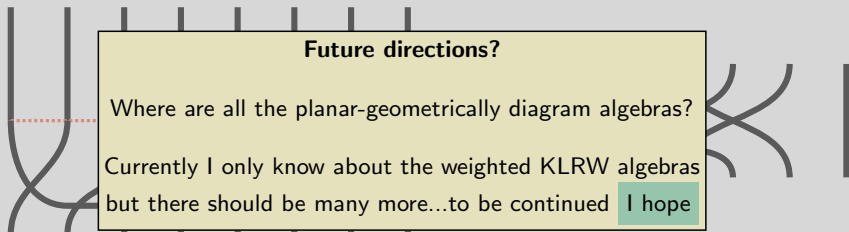
Define a diagram algebra that uses the distance in  $\mathbb{R}^2$

The result is called **weighted KLRW algebra**

These are “planar-geometrically symmetric group diagram algebras”

## String diagrams – the baby case

- ▶ We clearly have  $g(hf) = (gh)f$
- ▶ There is a do nothing operation  $1g = g = g1$



**Future directions?**

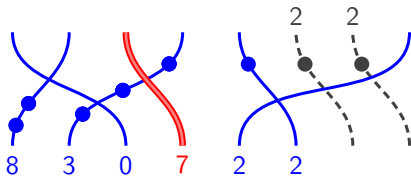
Where are all the planar-geometrically diagram algebras?

Currently I only know about the weighted KLRW algebras but there should be many more...to be continued I hope

- ▶ Generators–relations (the Reidemeister moves)



# Weighted string diagrams



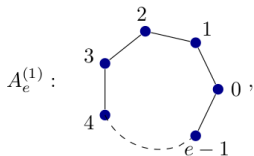
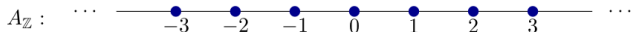
- ▶ Strings come in three types, **solid**, **ghost** and **red**

$$\text{solid} : \begin{array}{|c} \hline \\ \hline \end{array}, \quad \text{ghost} : \begin{array}{|c} \hline \cdots \\ \hline \end{array}, \quad \text{red} : \begin{array}{|c} \hline \\ \hline \end{array},$$

*i* *i* *i*

- ▶ Strings are labeled, and solid and ghost strings can carry dots
- ▶ Red strings **anchor** the diagram (red strings  $\leftrightarrow$  level)
- ▶ Otherwise no difference to symmetric group diagrams

# Weighted string diagrams



Examples of quivers  $\Gamma$



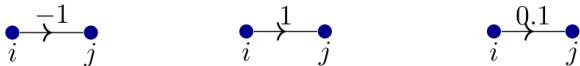
- ▶ The strings are labeled by  $i \in I$  from a fixed quiver  $\Gamma = (I, E)$
- ▶ The relations (that I am not going to show you ;-)) depend on  $e \in E$ , e.g.:

$$\begin{array}{c} \vdots \\ \bullet \\ \vdots \end{array} \begin{array}{c} | \\ | \\ | \end{array} = \begin{array}{c} \text{---} \\ \diagdown \text{---} \\ \diagup \text{---} \\ \text{---} \end{array} + \begin{array}{c} \vdots \\ \bullet \\ \vdots \end{array} \begin{array}{c} | \\ | \\ | \end{array} \quad \text{if } i \rightarrow j$$

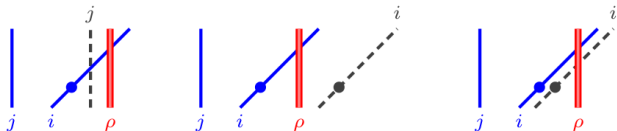
# Weighted string diagrams

$$X = (-2\sqrt{3}, -\sqrt{2}, 0.5, \pi, 5) \iff \begin{array}{cccccc} | & | & || & | & | & | \\ -2\sqrt{3} & -\sqrt{2} & 0 & 0.5 & \pi & 5 \end{array}$$

Weighted quiver



diagram



- ▶ Choose endpoints  $\mathbf{x} = (x_1, \dots, x_n) \in \mathbb{R}^n$ ,  $\rho \in \mathbb{R}^\ell$  for the solid and red strings
- ▶ Choose a weighting  $\sigma: E \rightarrow \mathbb{R}_{\neq 0}$  of the underlying graph  $\Gamma = (I, E)$
- ▶ The weighted KLRW algebra **crucially depends** on these choices of endpoints! This is very different from “usual diagram algebras”

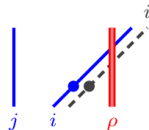
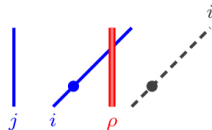
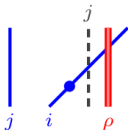
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Weighted quiver



diagram



## Weighting = ghost shifts

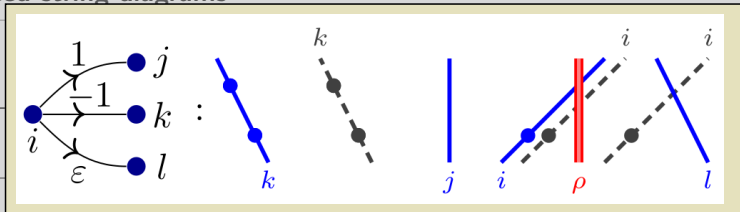
For  $\epsilon: i \rightarrow j, \sigma_\epsilon > 0$ , all solid  $i$ -strings get a ghost shifted  $|\sigma_\epsilon|$  units and mimicking it  
 For  $\epsilon: i \rightarrow j, \sigma_\epsilon < 0$ , all solid  $j$ -strings get a ghost shifted  $|\sigma_\epsilon|$  units and mimicking it

► The weight endpoints

This "asymmetric" definition, always shifting rightwards makes life a bit more convenient

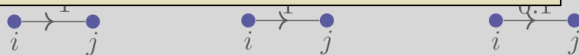
# Weighted string diagrams

$X = (-$

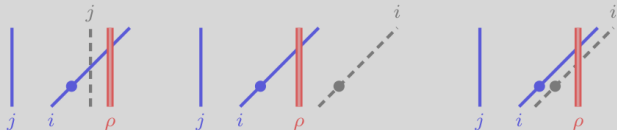


5

Weighted quiver



diagram



## Weighting = ghost shifts

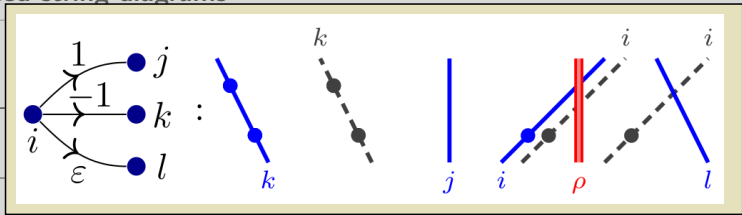
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- The weighted KLRW algebra **crucially depends** on these choices of endpoints! This is very different from “usual diagram algebras”



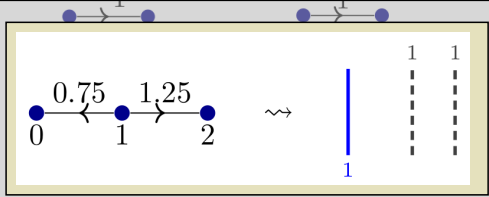
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$X = (-$

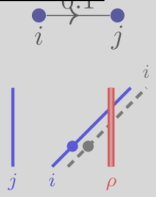


5

Weighted quiver



diagram



### Weighting = ghost shifts

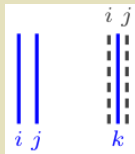
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## Weighted string diagrams

$$X = (-2\sqrt{3}, -\sqrt{2}, 0.5, \pi, 5) \leftrightarrow \begin{array}{cccccccc} | & & | & & || & | & & | & & | \\ \hline & & & & \circ & \circ & & & & \end{array}$$

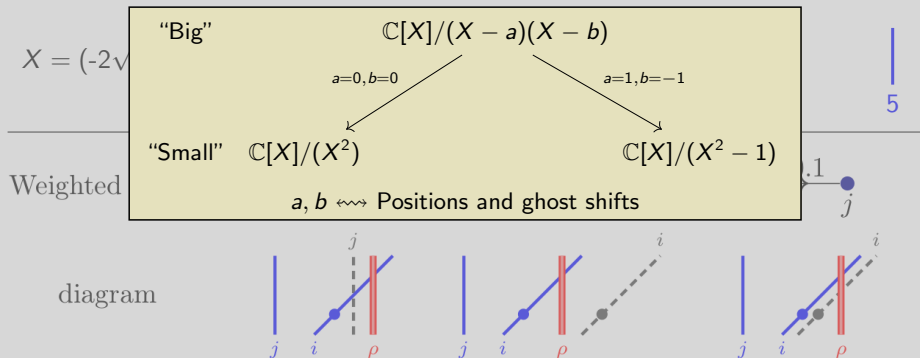
The following  $i$  and  $j$ -strings are not close:



**Slogan** Ghosts prevent the diagrams from being scale-able as for “usual diagram algebras”

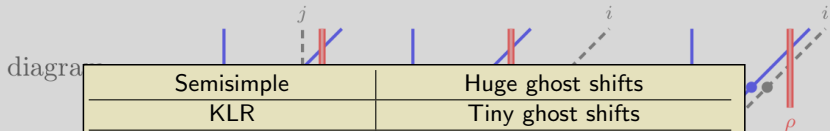
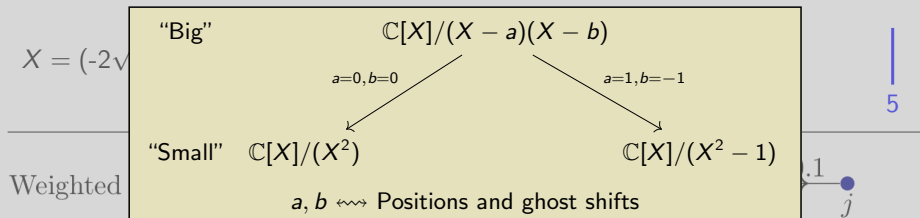
- ▶ Choose endpoints  $\mathbf{x} = (x_1, \dots, x_n) \in \mathbb{R}^n$ ,  $\rho \in \mathbb{R}^\ell$  for the solid and red strings
- ▶ Choose a weighting  $\sigma: E \rightarrow \mathbb{R}_{\neq 0}$  of the underlying graph  $\Gamma = (I, E)$
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# Weighted string diagrams



- ▶ Choose endpoints  $\mathbf{x} = (x_1, \dots, x_n) \in \mathbb{R}^n$ ,  $\rho \in \mathbb{R}^\ell$  for the solid and red strings
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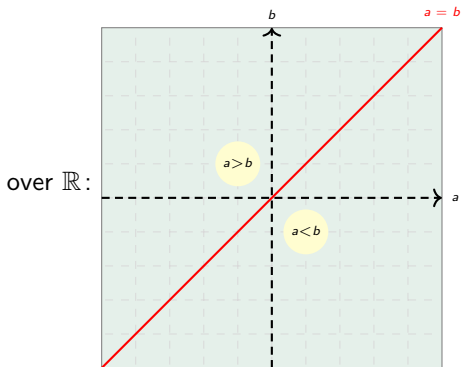
# Weighted string diagrams



Semisimple	Huge ghost shifts
KLR	Tiny ghost shifts
Quiver Schur	Some specific “cluster” spacing
Diagrammatic Cherednik	Ghost shifts 1
Unnamed algebras	The rest

- ▶ Choose endpoints  $x = (x_1, \dots, x_n) \in \mathbb{C}^n$ ,  $p \in \mathbb{C}$  for the solid and red strings
- ▶ Choose a weighting  $\sigma: E \rightarrow \mathbb{R}_{\neq 0}$  of the underlying graph  $\Gamma = (I, E)$
- ▶ The weighted KLRW algebra **crucially depends** on these choices of endpoints! This is very different from “usual diagram algebras”

# Hyperplanes



- ▶  $\mathbb{C}[X]/(X - a)(X - b)$  comes in two isomorphism classes:

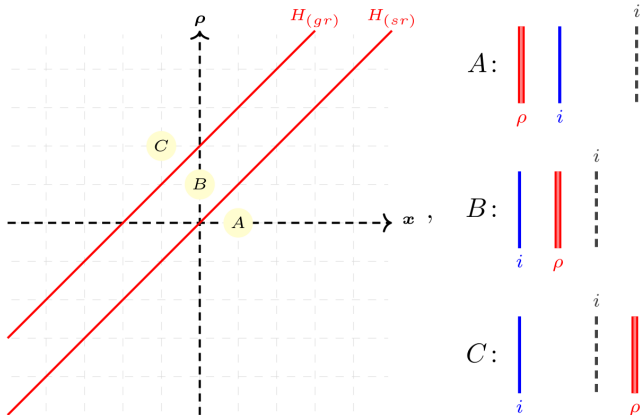
one double root  $a = b$

&

two different roots  $a \neq b$

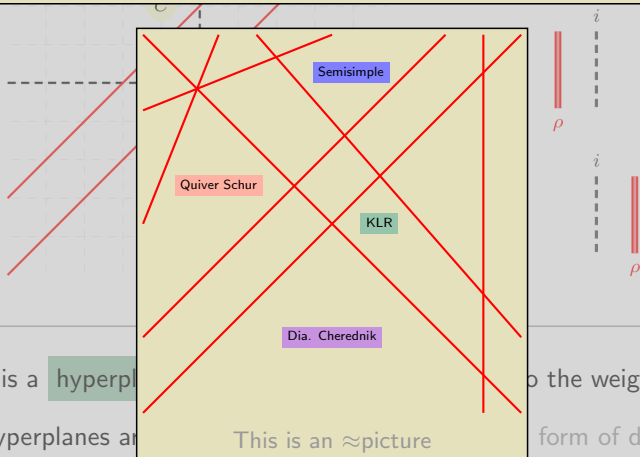
- ▶ What is the analog picture for weighted KLRW algebras?

# Hyperplanes



- ▶ There is a **hyperplane arrangement (HA)** associated to the weighted KLRW
- ▶ The hyperplanes are defined by **“colliding strings”** (a form of distance)

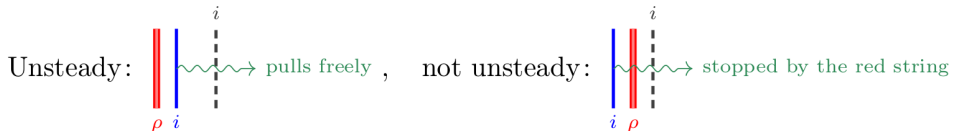
- Hy ▶ Alcoves of the HA  $\Rightarrow$  Morita equivalence classes of weighted KLRW algebras
- ▶ There is a theory of translation functors
  - ▶  $\approx$ picture 1 There is an alcove for KLR, an alcove for the semisimple case etc.
  - ▶  $\approx$ picture 2 Translation functors interpolate between these algebras



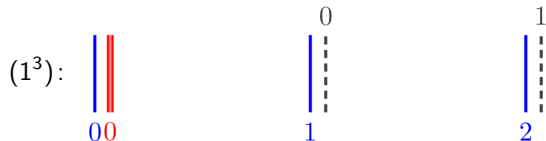
- ▶ There is a hyperplane to the weighted KLRW
- ▶ The hyperplanes are (in form of distance)
- This is an  $\approx$ picture

# Distance is it!

- ▶ Cyclotomic quotients  $\Leftrightarrow$  bounded regions:



- ▶ Cellular bases  $\Leftrightarrow$  minimal regions (I will elaborate momentarily):

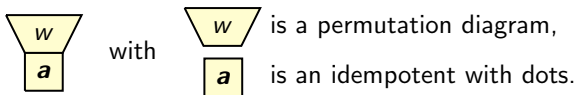


- ▶ More properties I won't explain today due to time restrictions...

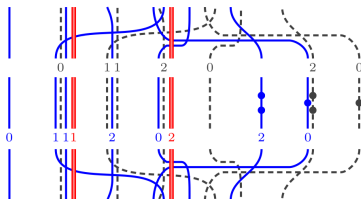
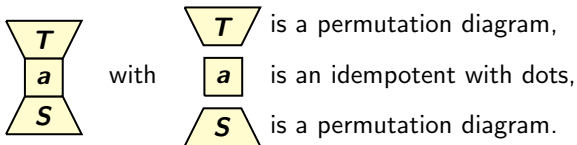


## Distance is it!

- ▶ Weighted KLRW algebras have **standard bases**, with the picture:



- ▶ Weighted KLRW algebras have **“cellular” bases**, with the picture:



Distance is

▶ Standard bases work regardless of the quiver but have no other property despite being a basis

▶ Weight

▶ Cellular bases depend on the quiver and give a classification of simple modules

(▶ Strictly speaking I should write “affine or sandwich” cellular but let us ignore that)

▶ Weighted KLRW algebras have “cellular” bases, with the picture:



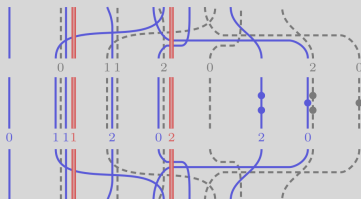
with



is a permutation diagram,

is an idempotent with dots,

is a permutation diagram.



Distance is

▶ Weight

▶ Standard bases work regardless of the quiver but have no other property despite being a basis

▶ Cellular bases depend on the quiver and give a classification of simple modules

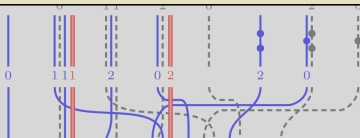
(▶ Strictly speaking I should write “affine or sandwich” cellular but let us ignore that)

▶ The overall strategy to construct cellular bases is the same for all type (but the details differ)

and for the infinite dimensional and the cyclotomic case the construction is also the same

▶ We know that the cellular bases work in types  $A, B, C$ , affine  $A, C, A^{(2)}, D^{(2)}$  but they should work even more general

▶ The combinatorics is inspired by, but different from, constructions of **Bowman**  $\sim 2017$ , **Ariki–Park**  $\sim 2012/2013$ , **Ariki–Park–Speyer**  $\sim 2017$



Distance is

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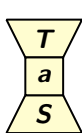
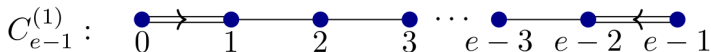
▶ The combinatorics is inspired by, but different from, constructions of **Bowman** ~2017, **Ariki–Park** ~2012/2013, **Ariki–Park–Speyer** ~2017

I will now indicate how the construction works in type  $C_{e-1}^{(1)}$

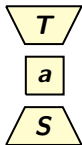
Why this type?

Because the code I am going to use works best for this type ;-)

# Minimal diagrams in type $C_{e-1}^{(1)}$



with

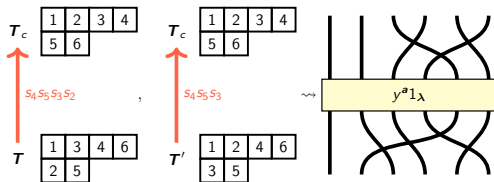


is a permutation diagram,

is an idempotent with dots,

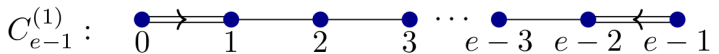
is a permutation diagram.

- ▶ The definition of the permutation follows the usual strategy in this context:



- ▶ Let me focus on the middle  $y^a 1_\lambda$

# Minimal diagrams in type $C_{e-1}^{(1)}$

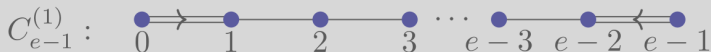


$C_3^{(1)} : (12, 6^3, 5) \leftrightarrow$

0	1	2	3	2	1	0	1	2	3	2	1
1	0	1	2	3	2						
2	1	0	1	2	3						
3	2	1	0	1	2						
2	3	2	1	0							

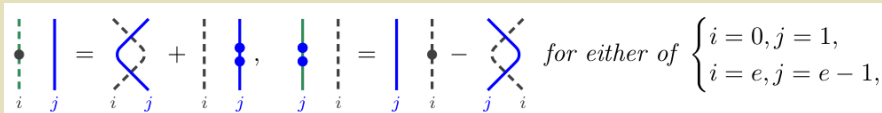
- ▶ Assume the tableaux combinatorics is given
- ▶ Place strings inductively as far to the right as possible (this is the order!)
- ▶  $1_\lambda$  is minimal with respect to placing the strings to the right
- ▶  $1_\lambda$  stays minimal when dots are put on certain strands  $\rightsquigarrow$  get  $y^a 1_\lambda$
- ▶ Done!

# Minimal diagrams in type $C_{e-1}^{(1)}$



0	1	2	3	2	1	0	1	2	3	2	1
1	0	1	2	3	2						

Lets ignore the dots for today – I bothered you with too much combinatorics anyway ;-)  
 But they come directly from the Reidemeister II relations, e.g.



In other words: **Stare at Reidemeister II !**


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# Minimal diagrams in type $C_{e-1}^{(1)}$


$C_{e-1}^{(1)}$  :


$C_3^{(1)} : (12)$

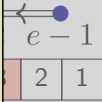
**Example for the middles  $y^a 1_\lambda$**

$y_{(1^5)} \mathbf{1}_{(1^5)} =$  

$y_{(5)} \mathbf{1}_{(5)} =$  

$y_{(2,1^3)} \mathbf{1}_{(2,1^3)} =$  

$y_{(4,1)} \mathbf{1}_{(4,1)} =$  



- ▶ Assume the t
- ▶ Place strings
- ▶  $1_\lambda$  is **minimal**
- ▶  $1_\lambda$  stays **min**
- ▶ **Done!**

is the order!)

ht  
get  $y^a 1_\lambda$



## Minimal diagrams in type $C_{e-1}^{(1)}$



### Wrap up

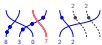
- ▶ Weighted KLRW algebras generalize KLR algebras and friends
  - ▶ They have a built-in distance
  - ▶ Most properties can be described using distance
  - ▶ Most properties are type-independent
  - ▶ Some properties should be (in some form) type-independent (order!)
- ▶  $1_\lambda$  is minimal with respect to placing the strings to the right
- ▶  $1_\lambda$  stays minimal when dots are put on certain strands  $\rightsquigarrow$  get  $y^a 1_\lambda$
- ▶ Done!

Where are we?



- Khovanov-Lauda-Rouquier – 2008 + many others (including OHT)
- KLR algebras are at the heart of categorical representation theory
- Similarly for quiver Schur algebras and diagrammatic Chevalky algebras
- **Problem:** All of these are actually really complicated!

Weighted string diagrams



- Strings come in three types, **solid**, **ghost** and **red**
- Solid strings are labeled, and solid and ghost strings can carry dots
- Red strings **anchor** the diagram (red strings  $\rightarrow$  level)
- Otherwise no difference to symmetric group diagrams

► Alluces of the HA  $\rightarrow$  Morita equivalence classes of weighted KLRW algebras

► There is a theory of translation functors

► **Hypercrite 1:** There is an algebra for KLR, an algebra for the semisimple case etc.

► **Hypercrite 2:** Translation functors interpolate between these algebras

► There is a **hyperplane** (this is an equivalence form of distance)

► The hyperplanes

Where are we?

Observation

It often helps to find a "bigger" interpolating algebra, e.g.:

"Big"  $\mathbb{C}[X]/(X - a)(X - b)$

"Small"  $\mathbb{C}[X]/(X^2)$

How to play the interpolation game using **plectic geometry**?

As an upshot: we get an algebra interpolating between various algebras appearing in categorical representation theory

► KLR

► SL

► The takeaway keyword: **Distance**

► **Problem:** All of these are actually really complicated!

Weighted string diagrams

$x = (-2\sqrt{3}, -\sqrt{2}, 0.5, \pi, 5) \rightarrow$

Weighted quiver

diagram

- Choose endpoints  $a = \{a_1, \dots, a_n\} \subset \mathbb{R}^n$ ,  $p \in \mathbb{R}^n$
- Choose a weighting  $\alpha: E \rightarrow \mathbb{R}_{\geq 0}$  of the underlying graph  $F = (V, E)$
- The weighted KLRW algebra **locally diagonal** on those choices of endpoints! This is very different from "usual diagram algebra"

Distance is it!

- Weighted KLRW algebras have **standard bases**, with the picture:
- Weighted KLRW algebras have **"cellular" bases**, with the picture:

String diagrams – the baby case

- We clearly have  $g(hf) = (gh)f$
- There is a **do nothing operation**  $1g = g = g1$

Generators-relations (the Reidemeister moves)

gens:  $\times$ ,  $\circ$ ,  $\otimes$

Weighted string diagrams

$x = i$

Weighted quiver

diagram

Weighting = ghost shifts

For  $i, j \rightarrow j, \sigma_i > 0$ , all solid  $\rightarrow$ strings get a ghost shifted  $\rightarrow$  units and mixing in the  $i, j, \sigma_i > 0$ , all solid  $\rightarrow$ strings get a ghost shifted  $\rightarrow$  units and mixing in the endpoints! This is very different from "usual diagram algebra"

Minimal diagrams in type  $C_{n-1}^{\pm}$

Example for the minimal  $y^* 1_{\pm}$

► Assume the  $1_{\pm}$  is in **matrix** (in the order!)

►  $1_{\pm}$  **ways out** get  $y^* 1_{\pm}$

► **Done!**

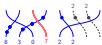
There is still much to do...

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► Alluces of the HA  $\rightarrow$  Morita equivalence classes of weighted KLRW algebras

► There is a theory of translation functors

► **Hypercube 1:** There is an alcove for KLR, an alcove for the semisimple case etc.

► **Hypercube 2:** Translation functors interpolate between these alcoves

► There is a **hyperplane** (this is an equidistant form of distance)

► The hyperplanes

Where are we?

Observation

It often helps to find a "bigger" interpolating algebra, e.g.:

"Big"  $\mathbb{C}[X]/(X - a)(X - b)$

"Small"  $\mathbb{C}[X]/(X^2)$   $\xrightarrow{a=b=0}$   $\mathbb{C}[X]/(X^2 - 1)$

How to play the interpolation game using **plectic geometry**?

► KLR

► SL

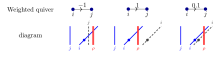
► The tabular keyword: **Distance**

► **Problem:** All of these are actually really complicated!

Weighted string diagrams

This is the first time I use the number  $\tau$  (is a  $\tau$  is a  $\tau$ )

$$X = (-2\sqrt{3}, -\sqrt{2}, 0.5, \tau, 5) \rightarrow \begin{array}{c} \tau \\ \tau \\ \tau \\ \tau \\ \tau \end{array}$$



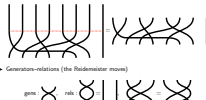
- Choose endpoints  $a = \{a_1, \dots, a_n\} \in \mathbb{R}^n$ ,  $p \in \mathbb{R}^n$  for the solid and red strings
- Choose a weighting  $\alpha: E \rightarrow \mathbb{R}_{\geq 0}$  of the underlying graph  $F = (I, E)$
- The weighted KLRW algebra is **locally diagrammatic** on these choices of endpoints! This is very different from "usual diagram algebra"

Distance is it!

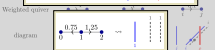
- Weighted KLRW algebras have **standard bases**, with the picture:
 
 with  $\begin{array}{c} \tau \\ \tau \\ \tau \\ \tau \\ \tau \end{array}$  is an idempotent with dots.
- Weighted KLRW algebras have **"cellular" bases**, with the picture:
 
 with  $\begin{array}{c} \tau \\ \tau \\ \tau \\ \tau \\ \tau \end{array}$  is a permutation diagram.

String diagrams – the baby case

- We clearly have  $g(hf) = (gh)f$
- There is a do nothing operation  $1g = g = g1$



Weighted string diagrams



- Weighting = ghost shifts
- For  $i = j \rightarrow j, \alpha_i > 0$ , all solid  $\rightarrow$ strings get a ghost shifted  $\tau$  units and missing in the  $i = j, \alpha_i < 0$ , all solid  $\rightarrow$ strings get a ghost shifted  $\tau$  units and missing in the  $i = j$
- The weighted KLRW algebra is **locally diagrammatic** on these choices of endpoints! This is very different from "usual diagram algebra"

Minimal diagrams in type  $C_{n-1}^{\tau}$

Example for the minimal  $y^* \tau_{1,2}$

$C_{n-1}^{\tau}$

$C_{n-1}^{\tau}(\tau)$

► Assume the  $\tau_{1,2}$  is **minimal**

► Place strings

►  $\tau_{1,2}$  is **minimal**

►  $\tau_{1,2}$  **always dot**

► **Done!**

Thanks for your attention!