

Everything I do today works in general type, but mostly focus on S_n for simplicity.

Famously, S_n has Coxeter presentation: gens $S_i = (i \ i+1)$ for $i \in \{1, \dots, n-1\}$ (2)

$S_n = \langle S_i \rangle / S_i^2 = 1$ quadratic ← also a presentation for $[S_n]$

$S_i S_{i+1} S_i = S_{i+1} S_i S_{i+1}$
 $S_i S_j = S_j S_i \quad (j \neq i \pm 1)$ } braid

For my story it helps to discuss some variants on $[S_n]$ as well, keep braid rels but modify quadratic

*-monoid

$\langle S_i \rangle / S_i^2 = S_i$ braid

n) Coxeter algebra

$N(S_n) = \langle D_i \rangle / D_i^2 = 0$ braid

Jucys-Murphy algebra

ALL VARIANTS FROM / GEOMETRY OF FLAGS!

Monoid algebra is "O-Hecke algebra"

All these have a basis $\leftrightarrow S_n$, defined as follows: (focus on NC)

$\underline{w} = (i_1, i_2, \dots, i_d)$ a word in S_n , what I call an expression.

\underline{w} is reduced expression or rex for w if $w = S_{i_1} \dots S_{i_d}$ and no shorter word suffices, $l(w) = d$.

$D_{\underline{w}} = D_{i_1} D_{i_2} \dots D_{i_d}$ for any expr.

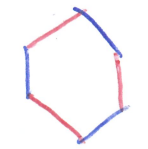
Def: $D_{\underline{w}} = D_{\underline{w'}}$ when \underline{w} a rex for w .

Well-defined (indep of \underline{w}) b/c:

Thm (Matsumoto): Any two rexes for w connected by braid rels.

All variants have "some behavior" for rexes. Differ in how to interpret non-rexes.

GO TO PICTURE AND DISCUSS WALKS AROUND



NC has nice feature of isolating rexes as special:

$D_{\underline{w}} \neq 0 \iff \underline{w}$ is a rex.

We'll generalize v to a category w/ objects $I \subset S$ (each variant) to which we associate $W_I = \langle s_i \mid i \in I \rangle$ ③

Ex: In S_n , $W_I \cong S_{n_1} \times \dots \times S_{n_k}$ where $\sum n_k = n$.

In general we only consider I with W_I finite \implies has largest element w_I .

In our category $\text{Hom}(I, J)$ has basis $\leftrightarrow \frac{W_J}{W_I}$ (should probably reverse because composition R to L annoying and I may forget for pedagogy...)

$\text{Hom}(\phi, \phi) =$ previous algebra

composition $(\phi \circ s \leftarrow \phi) \leftrightarrow$ generator Generator: $I \leftrightarrow I \cup \{s\} =: I_s$

shorten $[\phi, s, \phi]$

this composition is never invertible. Think: $[s, \phi]$ is not a permutation but a map from $\{1, \dots, n\}$ to $\{1, 2, 3, \dots, n\}$

Can do for $\mathbb{C}[S_n]$ but need to change generators, will instead focus on NC and $*$.

- Need to
- 1) Motivate
 - 2) Describe
 - 3) State awesome theorems

2) Is easy in $*$ -monad variant, see. We'll say how generator acts on basis of double cosets. If $p \in \frac{W_I}{W_J}$ then $\exists!$ $q \in \frac{W_{I_s}}{W_J}$

with $p \circ q$: $[I_s, I] * p = q$.

Conversely, q splits into ≥ 1 cosets for $\frac{W_I}{W_J}$, and one is maximal in bracket order. E.g. $\bar{q} \circ q$ is ! max'l element, then $p \circ q$ maximal

if $\bar{p} = \bar{q}$.

$[I, I_s] * q = p$.

SHOW OFF IN PICTURE.

Aside! Easiest way to describe mult. of elements is $SC \cong \text{Kar}(*\text{-monad})$
 Idempotents = $\{w_I\}_{I \subset S}$

~~Def:~~ Def: A double cost expression $\underline{I} = [I_0, I_1, \dots, I_d]$ is (4)

a composition of generators, i.e. $I_{i+1} = I_i s$ or $I_i = I_{i+1} s$ for some s .

Can assign a length, $l(\underline{I}, I_r) = l(\underline{I} s, I_r) = l(I_r) - l(\underline{I})$ where

$l(I) = l_W(w_I)$. DISCUSS LENGTHS ON PIC.

Can assign a length to double cost, $l(p) = 2 l_W(p) - l(I) - l(J)$

$p \in \{w\} \subseteq \frac{W}{W_I} / \frac{W}{W_J}$

Note! For $p \in \{w\} \subseteq \frac{W}{W_I} / \frac{W}{W_J}$, $l(p) = 2 l_W(w)$.

Def: \underline{I} is reduced if express p with $l(\underline{I}) = l(p)$.

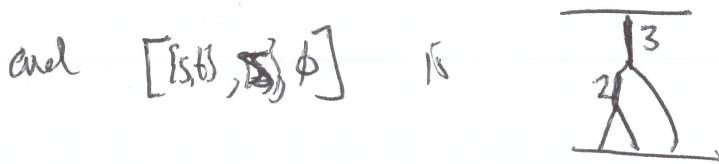
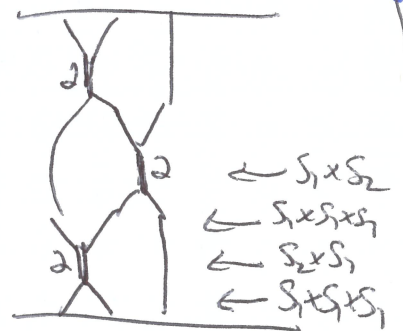
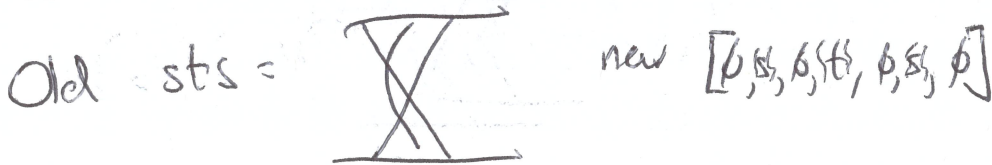
(One of many equiv. defns.)

In picture, reflex are oriented paths from base.

base = $\frac{W}{W_I} / \frac{W}{W_I}$ cost containing 1. $l(\text{base}) = 2 l_W(w_I) - l(I) - l(I) = 0$.

... is not reduced in either direction. (Maybe do root hyperplanes here!)

For those who like webs, in type A can draw expressions.

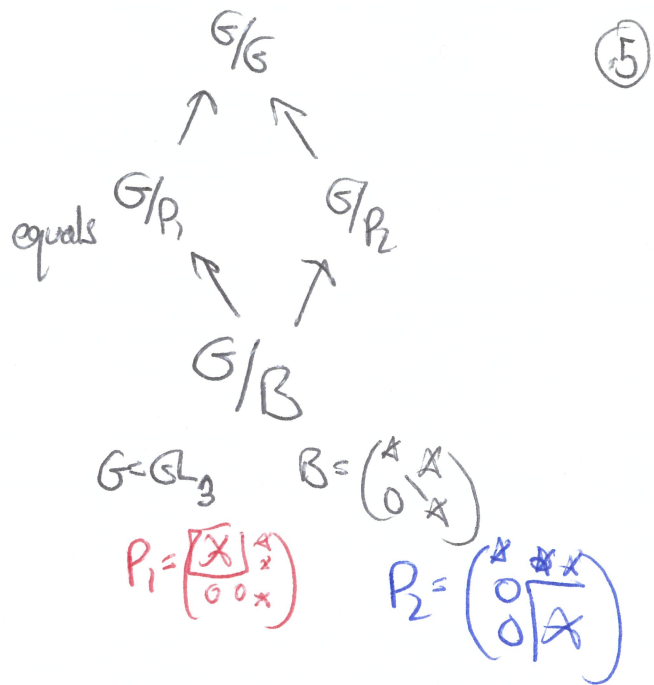
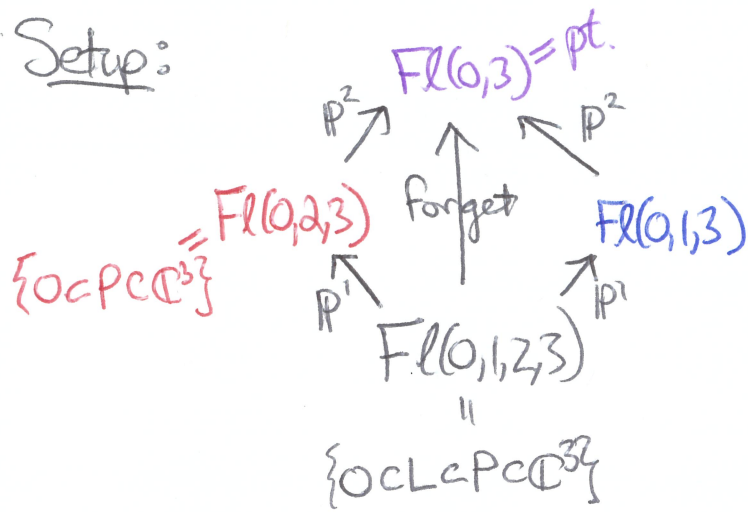


Geom/Alg source

Easiest to explain using NC first.

Note! In k -version, this category is Karasik envelope of k -normal I -degenerations \rightarrow W_I

Setup:



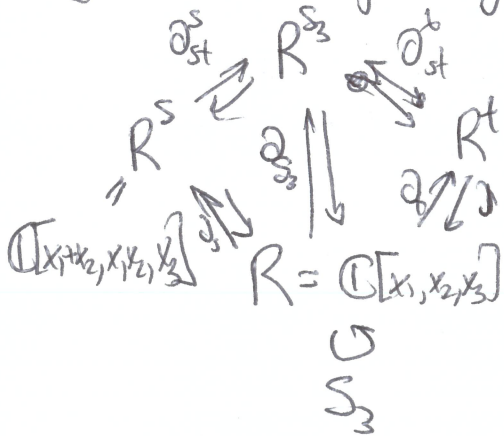
To this setting one can apply all sorts of functors, e.g. $H^*(-)$.



- 1) Draw pullbacks, they're inclusions, ring homoms.
- 2) Discs performed below.
- 3) Draw pushforwards, they're surjections, NOT ring hom.
- 4) Diagram commutes.

When $Y \xrightarrow{P^k} X$ have $H^*(Y) \xrightarrow{\pi_{k*}} H^{*-2k}(X)$ "integrate along fibers"

Can do this algebraically:



(to get above picture kill ideal $(R_+^{S_3})$.
(this is really for (equiv cohom.))

The maps ∂_s are Denazire operators.

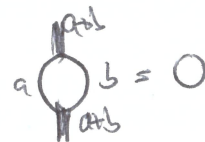
$$\partial_s(f) = \frac{f - sf}{x_1 - x_2}, \text{ den } \text{divisor num.}$$

Ex: $\partial(x_1^3) = \frac{x_1^3 - x_2^3}{x_1 - x_2} = x_1^2 + x_1x_2 + x_2^2$

Thm (EK) (se)

1) $NC(\mathbb{Z})$ has presentation w/ relns

a) "quadratic" $[I_S, I, I_S] = 0$
 for $\mathbb{Z}([I_S, I, I_S] = 0)$



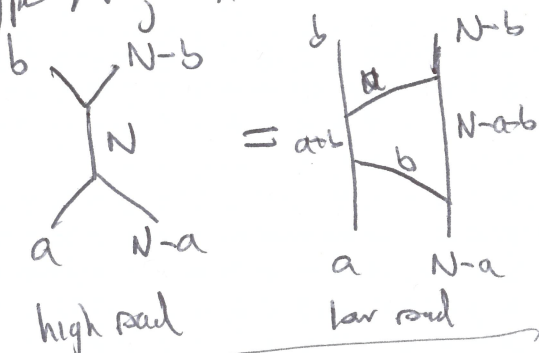
braided

b) $[I_{S^*}, I_S, I] = [I_{S^*}, I_S, I]$
 and reverse



c) switchback. One such relation for each
 imed finite Cox gp and each pair S, S
 with $t \neq w_0 S w_0$

In type A_N if $a+b < N$ then



general idea
 Both reflex for $W_1/S_N/W_2$
 coset containing w_0

and horiz flip.

Together with distinct switchbacks

like $\Delta Y = \Delta Y$

2) $\partial_I = 0 \iff \underline{I}$ is not reduced

3) "Matsumoto": All reflex for p related by braid relations
 (so ∂_p defined)

4) $\{\partial_p\}_{p \in W_1/W_2}$ a basis for $\text{Hom}_{\mathbb{Z}}(R^J, R^I)$

5) $\prod_{n \geq 0} NC(S_n) \cong \text{AssGr}(\text{Type A Webs})$
 for glas

6) Bruhat order via "subexpressions"

Remarks: 1) Defn of rex equivalent to Williamson's but much less technical.

(8)

2) $[I_0 c \dots c k_1 \dots \rightarrow I_1 c \dots c k_2 \dots \rightarrow I_2 c \dots]$ is reduced iff
lengths add in $w_{k_1} (w_{I_1}^{-1} w_{k_2}) (w_{I_2}^{-1} w_{k_3}) \dots$

3) Higher braid rels (egs 3-cells, etc) not known.

4) In any type get switchback via switchback leapfrog

$$u_1 = w_{s_1} u_0 \quad u_0 = s \quad u_1 = w_{s_3} u_1 w_{s_5} \quad u_2 = w_{s_1} s w_{s_1} \quad u_3 = w_{s_2} u_1 w_{s_2} \dots$$



$$[s, s, t] = [\hat{u}_0, \hat{u}_0 \hat{u}_1, \hat{u}_1, \hat{u}_1 \hat{u}_2, \hat{u}_2, \hat{u}_2 \hat{u}_3, \hat{u}_3, \hat{u}_3 \hat{u}_4, \hat{u}_4, \hat{u}_4 \hat{u}_5, \hat{u}_5]$$

How far do you go? It depends - could be longer or shorter than period!
But ends in \hat{t} .

Applications: 1) Braid rels ~~are~~ fundamental to Hecke category + light leaves basis.

KGLP: light leaves for singular Soergel bimodules.

2) Specifically, $\partial [s, s, t] = \partial \underline{\mathbb{I}} \iff R^s \underset{R^s}{\otimes} R^t \overset{\oplus}{\subset} BS(\underline{\mathbb{I}})$

3) Presentation of Hecke algebra should come out of this easily.

4) Rex combinatorics of a completely new type. High roads, low roads, atomic cases...

5) Likely related to Stanley's funny rex counting results.

2') (Williamson's original part) $BS(\underline{\mathbb{I}})_{\text{var}}$ as resolution of sing. of $\overline{\mathcal{O}}_p$.

Things that didn't make it into this talk

(19)

1) If $t = w_0 s w_0$ then $[S, S, t]$ is unique rex. for $\bar{p} = w_0$

2) Redundancy subgps. P , Howlett.

3) Using PD to define pushforward.

4) $P_I / S / P_J \iff W_I / W / W_J$ coset bruhat.

5) Subexpressions as stralls

6) Reduced via idea of ~~root~~ hyperplanes

Never do

